

The class of $(2P_3, C_4, C_6)$ -free graphs, part I: $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0

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Abstract

This is the first in a series of two papers dealing with $(2P_3, C_4, C_6)$ -free graphs, or equivalently, $(2P_3, \text{even hole})$ -free graphs. In this two-paper series, we give a full structural description of $(2P_3, C_4, C_6)$ -free graphs that contain no simplicial vertices, and we show that such graphs have bounded clique-width. This implies that GRAPH COLORING can be solved in polynomial time for $(2P_3, C_4, C_6)$ -free graphs. In this paper, we describe the structure of $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 (where T_0 is a certain 2-connected graph on nine vertices in which all holes are of length five), and we show that such graphs either contain a simplicial vertex or have bounded clique-width. In the second part of this series, we describe the structure of all $(2P_3, C_4, C_6, C_7, T_0)$ -free graphs that contain no simplicial vertices, and we show that such graphs have bounded clique-width. The full statement of the theorem describing the structure of $(2P_3, C_4, C_6)$ -free graphs that contain no simplicial vertices is given in the second paper of this series.

1 Introduction

All graphs in this paper are finite, simple, and nonnull. We mostly use standard terminology and notation, formally introduced in section 2. For now, let us define a few basic terms. As usual, for a positive integer k , K_k is the complete graph on k vertices, P_k is the path on k vertices and $k - 1$ edges, and C_k ($k \geq 3$) is the cycle on k vertices. For a positive integer k and a graph H , the disjoint union of k copies of the graph H is denoted by kH . (See Figure 1.1 for some graphs relevant to this paper.) For a graph H , we say that a graph G is H -free if no induced subgraph of G is isomorphic to H . For a family of graphs $\mathcal{H}$, we say that a graph G is $\mathcal{H}$ -free if G is H -free for all $H \in \mathcal{H}$. A *hole* in a graph G is an induced cycle on at least four vertices, and the *length* of a hole is the number of vertices (equivalently: edges) that it contains. A hole is *even* (resp. *odd*) if its length is even (resp. odd). A *clique* (resp. *stable set*) is a (possibly empty) set of pairwise adjacent (resp. nonadjacent) vertices, and a *simplicial vertex* is a vertex whose neighbors form a (possibly empty) clique. A *proper coloring* of a graph G is an assignment of colors to the vertices of G in such a way that no two adjacent vertices receive the same color. For an integer k , a graph G is said to be k -colorable if there exists a proper coloring of G that uses at most k colors. The *chromatic number* of G , denoted by $\chi(G)$, is the smallest integer k such that G is k -colorable. GRAPH COLORING is the following problem.

GRAPH COLORING

Instance: A graph G and an integer k .

Question: Is G k -colorable?

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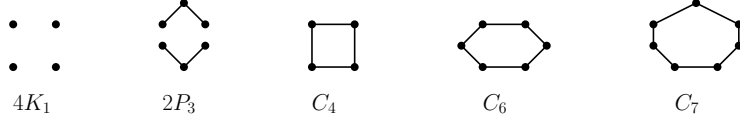


Figure 1.1: From left to right: graphs $4K_1$, $2P_3$, C_4 , C_6 , and C_7 .

This paper is the first in a series of two papers that deal with $(2P_3, C_4, C_6)$ -free graphs. In this series, we give a full structural description of $(2P_3, C_4, C_6)$ -free graphs that contain no simplicial vertices, and we show that GRAPH COLORING can be solved in polynomial time for $(2P_3, C_4, C_6)$ -free graphs.

Before describing our results in more detail, let us make some general remarks and give the relevant context for our work. $(2P_3, C_4, C_6)$ -free graphs are trivially recognizable in $O(n^6)$ time. Further, the MAXIMUM CLIQUE problem can be solved in polynomial time for C_4 -free graphs [1, 12, 21, 28];¹ the MAXIMUM STABLE SET problem can be solved in polynomial time for all (banner, P_8)-free graphs [13], where the *banner* is the five-vertex graph consisting of a C_4 with a pendant edge;² and the 3-COLORING problem can be solved in polynomial time for sP_3 -free graphs for any fixed positive integer s [3]. Therefore, the MAXIMUM CLIQUE, MAXIMUM STABLE SET, and 3-COLORING problems are all solvable in polynomial time for $(2P_3, C_4, C_6)$ -free graphs. As stated above, one of the main goals of our two-paper series is to show that the GRAPH COLORING problem can be solved in polynomial time for $(2P_3, C_4, C_6)$ -free graphs.

We note that $4K_1$ is an induced subgraph of $2P_3$, and so $(2P_3, C_4, C_6)$ -free graphs form a proper superclass of the class of $(4K_1, C_4, C_6)$ -free graphs. The structure of $(4K_1, C_4, C_6)$ -free graphs is almost completely understood: a full structural description of $(4K_1, C_4, C_6)$ -free graphs that contain an induced C_7 was given in [14], and a full structural description of $(4K_1, C_4, C_6, C_7)$ -free graphs that contain no simplicial vertices was given in [22]. Moreover, all the following problems can be solved in $O(n^3)$ time for $(4K_1, C_4, C_6)$ -free graphs: recognition, MAXIMUM CLIQUE, and GRAPH COLORING [23].³ On the other hand, we note that every hole of length at least eight contains an induced $2P_3$, and consequently, all holes in a $(2P_3, C_4, C_6)$ -free graph are of length five or seven; thus, $(2P_3, C_4, C_6)$ -free graphs are even-hole-free, and in fact, $(2P_3, C_4, C_6)$ -free graphs are precisely the $(2P_3, \text{even hole})$ -free graphs. The class of even-hole-free graphs has received a great deal of attention over the past couple of decades. In particular, decomposition theorems have been proven for even-hole-free graphs [7, 10], such graphs can be recognized in polynomial time [5, 8], and they are χ -bounded by a linear function [6]. As mentioned above, the MAXIMUM CLIQUE problem can be solved in polynomial time for C_4 -free graphs, and consequently, for even-hole-free graphs as well. On the other hand, the time complexity of the MAXIMUM STABLE SET and GRAPH COLORING problems is unknown for even-hole-free graphs. Finally, let us mention that the structure of “ ℓ -holed graphs,” i.e. graphs in which all holes are of length ℓ , is well understood for each integer $\ell \geq 7$ (see [9]). As pointed out above, all holes in a $(2P_3, C_4, C_6)$ -free graph are of length five or seven, and consequently, all $(2P_3, C_4, C_6, C_7)$ -free graphs are 5-holed. So, while general 5-holed graphs are

¹Indeed, any C_4 -free graph has only $O(n^2)$ maximal cliques [1, 12], and if a graph G has K maximal cliques, they can all be found in $O(Kn^3)$ time [21, 28]. So, all maximal cliques of a C_4 -free graph can be found in $O(n^5)$ time, and it immediately follows that the MAXIMUM CLIQUE problem can be solved in $O(n^5)$ time for C_4 -free graphs.

²Since the banner contains an induced C_4 , and P_8 contains an induced $2P_3$, we see that all $(2P_3, C_4, C_6)$ -free graphs are, in particular, (banner, P_8)-free.

³Clearly, the stability number of any $4K_1$ -free graph is at most three, and so the MAXIMUM STABLE SET problem can trivially be solved in $O(n^3)$ for $4K_1$ -free graphs, and consequently, for $(4K_1, C_4, C_6)$ -free graphs as well.

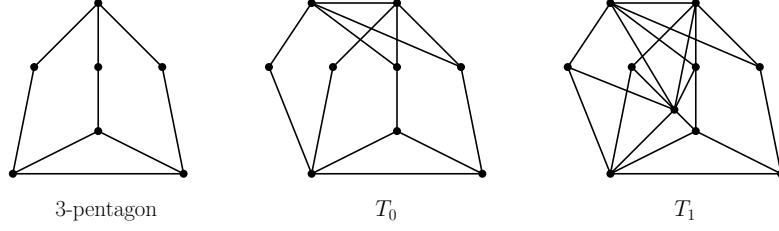


Figure 1.2: The 3-pentagon (left), T_0 (center), and T_1 (right).

not well understood, our results yield a good description of one particular subclass of 5-holed graphs, namely, that of $(2P_3, C_4, C_6, C_7)$ -free graphs.

Let us now describe our results in a bit more detail. Our main goal in this two-paper series is to give a full structural description of all $(2P_3, C_4, C_6)$ -free graphs that contain no simplicial vertices, and to prove that such graphs have bounded clique-width (for a formal definition of “clique width,” see section 6). Since simplicial vertices pose no obstacle to coloring in polynomial time,⁴ and since GRAPH COLORING can be solved in polynomial time for graphs of bounded clique-width [25], it follows that the GRAPH COLORING problem can be solved in polynomial time for $(2P_3, C_4, C_6)$ -free graphs. We note, however, that $(2P_3, C_4, C_6)$ -free graphs do not have bounded clique-width: only those $(2P_3, C_4, C_6)$ -free graphs that contain no simplicial vertices do. Indeed, it was shown in [2] that $4K_1$ -free chordal graphs have unbounded clique-width (a *chordal graph* is a graph that contains no holes, i.e. a graph in which all induced cycles are triangles), and clearly, $4K_1$ -free chordal graphs form a proper subclass of $(2P_3, C_4, C_6)$ -free graphs. However, it is well known that all chordal graphs have a simplicial vertex [11].

T_0 is the graph on nine vertices represented in Figure 1.2 (center); clearly, T_0 is 2-connected, and it is not hard to check that T_0 is $(2P_3, C_4, C_6, C_7)$ -free (for a formal proof, see Proposition 4.1), and in particular, all holes in T_0 are of length five. In this paper, we give a full structural description of $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 (see Theorem 5.1); we note that this structural description yields an $O(n^3)$ time recognition algorithm for $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 (see section 7). We also give a full structural description of all $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 , and contain no simplicial vertices (see Theorem 5.2), and we show that all such graphs have bounded clique-width (see Theorem 6.6). In the second part of this series [24], we give a full structural description of $(2P_3, C_4, C_6, C_7, T_0)$ -free graphs that contain no simplicial vertices, and we show that such graphs have bounded clique-width. Together, these two papers yield a full structural description of $(2P_3, C_4, C_6)$ -free graphs that contain no simplicial vertices, and the fact that all such graphs have bounded clique-width; as explained above, this implies that GRAPH COLORING can be solved in polynomial time for $(2P_3, C_4, C_6)$ -free graphs. The full formal statement of our structure theorem for $(2P_3, C_4, C_6)$ -free graphs that contain no simplicial vertices is given in the second paper of our series.⁵

It may be worth emphasizing that, while we were able to fully describe the structure of $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 (regardless of the pres-

⁴Indeed, if k is an integer, and x is a simplicial vertex of a graph G on at least two vertices, then G is k -colorable if and only if $d_G(x) \leq k - 1$ and $G \setminus x$ is k -colorable. Moreover, the simplicial vertices of a graph (if any) can easily be found by simply examining the neighborhood of each vertex and checking if that neighborhood is a clique.

⁵We omit the statement of this theorem in the present paper because, to make sense of the theorem, we would need some rather lengthy definitions that are relevant for $(2P_3, C_4, C_6, C_7, T_0)$ -free graphs, and those definitions are only given in the second paper.

ence or absence of simplicial vertices), in the $(2P_3, C_4, C_6, C_7, T_0)$ -free case, simplicial vertices complicate matters more substantially, and for this reason, in the second paper of our series, we restrict our attention to those $(2P_3, C_4, C_6, C_7, T_0)$ -free graphs that contain no simplicial vertices. Thus, even though we ultimately get a full structural description of $(2P_3, C_4, C_6)$ -free graphs that contain no simplicial vertices, this is not quite a full structure theorem for all $(2P_3, C_4, C_6)$ -free graphs. This is essentially because the class of $(2P_3, C_4, C_6)$ -free graphs is not closed under the addition of simplicial vertices, that is, adding a simplicial vertex to a $(2P_3, C_4, C_6)$ -free graph may possibly produce a graph that contains an induced $2P_3$,⁶ and is consequently not $(2P_3, C_4, C_6)$ -free. Of course, for many algorithmic problems (including GRAPH COLORING, but by no means limited to it), simplicial vertices can be handled with ease.

1.1 Further remarks on our results and some related open problems

Let us now briefly discuss some problems related to our work that still remain open. First, while GRAPH COLORING is indeed solvable in polynomial time for graphs of bounded clique-width, all known polynomial-time coloring algorithms for graphs of clique-width at most k have running time $O(n^{f(k)})$ for some fast growing function f (see [15] for an overview). So, it might be interesting to try to develop a faster coloring algorithm for $(2P_3, C_4, C_6)$ -free graphs, perhaps using the structural results described in our two-paper series. Despite substantial effort, we have thus far not been able to develop a coloring algorithm for this class that does not rely on clique-width. More precisely, using the results of [18, 19] and the structural results of the present paper, we can in fact color $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 in only $O(n^3)$ time (see section 7). (Here, we use the results of [18, 19] as a black box, but we do note that the algorithms in question rely on integer programming and are therefore not combinatorial.) However, in the second paper of our series, we are forced to rely on bounded clique-width to color $(2P_3, C_4, C_6, C_7, T_0)$ -free graphs, which yields a slow running time.

Second, note that $(2P_3, C_4, C_6)$ -free graphs form a proper subclass of the class of (P_7, C_4, C_6) -free graphs. It would, of course, be interesting to describe the structure of (P_7, C_4, C_6) -free graphs and to determine the time complexity of GRAPH COLORING for this class. Despite quite some effort, we have not succeeded in doing this. Nevertheless, in section 3, we in fact describe the structure of (P_7, C_4, C_6) -free graphs that contain an induced C_7 and do not admit a clique-cutset. This reduces the problem of describing the structure of (P_7, C_4, C_6) -free graphs to that of describing the structure of (P_7, C_4, C_6, C_7) -free graphs, as long as we are willing to regard the clique-cutset as an acceptable decomposition in this context.⁷ However, so far, we have not been able to make any meaningful progress on the structure of (P_7, C_4, C_6, C_7) -free graphs.

Here, it may be worth giving a bit more context. A *three-path-configuration* (or *3PC* for short) is any theta, pyramid, or prism (see Figure 1.3). A *wheel* is any graph that consists of a hole, plus an additional vertex that has at least three neighbors in the hole. A *Truemper configuration* is any 3PC or wheel. For a (slightly dated) survey on Truemper configurations, see [29]. Clearly, all holes in a (P_7, C_4, C_6, C_7) -free graph are of length five. Meanwhile, the only 3PC in which all holes are of length five is the 3-pentagon (see Figure 1.2, and note

⁶This is because $2P_3$ itself contains a simplicial vertex. On the other hand, since holes do not contain simplicial vertices, it is not possible to generate a new hole by simply adding one or more simplicial vertices, and in particular, any (C_4, C_6) -free graph remains (C_4, C_6) -free if we add a simplicial vertex to it.

⁷Note, however, that P_7 admits a clique-cutset, and consequently, the class of (P_7, C_4, C_6) -free graphs is not closed under the operation of gluing along a clique (the operation that “reverses” clique-cutsets). The clique-cutset decomposition is nevertheless reasonable to accept in this context, since there are powerful algorithmic tools for handling this decomposition (see [27]).

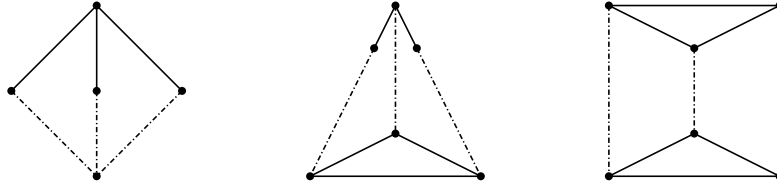


Figure 1.3: Three-path-configurations (3PCs): theta (left), pyramid (middle), and prism (right). A full line represents an edge, and a dashed line represents a path that has at least one edge.

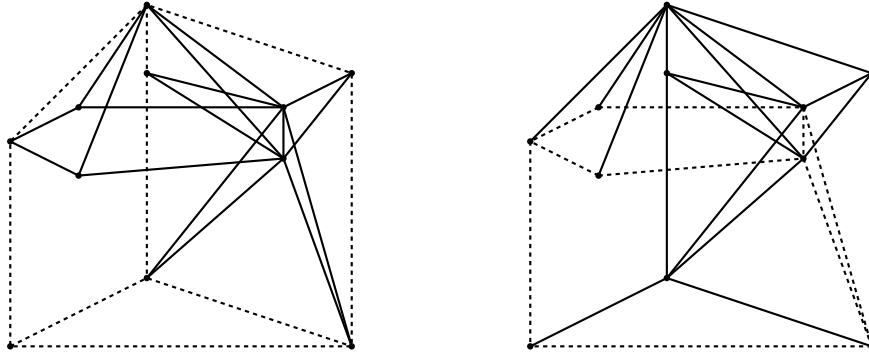


Figure 1.4: Two copies of the same (P_7, C_4, C_6, C_7) -free graph on eleven vertices, with induced 3-pentagons intersecting in an inconvenient way. Both full and dashed lines represent edges. The edges of two induced 3-pentagons are indicated by dashed lines.



Figure 1.5: The only wheels in which all holes are of length five.

that the 3-pentagon is a type of pyramid). Unfortunately, 3-pentagons may intersect in rather inconvenient ways in (P_7, C_4, C_6, C_7) -free graphs (see Figure 1.4 for an example), and this is the main reason why we have not been able to make meaningful progress on (P_7, C_4, C_6, C_7) -free graphs. Meanwhile, induced 3-pentagons in $(2P_3, C_4, C_6, C_7)$ -free graphs behave in a more manageable fashion (this is studied in the second paper of our series), which is why we were able to get a good structural description of such graphs, as long as they do not contain simplicial vertices. Finally, let us mention in passing that the two wheels represented in Figure 1.5 are the only wheels in which all holes are of length five; so, those two wheels are the only wheels that a (P_7, C_4, C_6, C_7) -free graph may possibly contain as induced subgraphs.

1.2 Paper outline

We complete the introduction with an outline of the paper. In section 2, we introduce the terminology and notation that we will use throughout the paper (see subsection 2.1), and we also prove a few simple results that will be of use to us as we proceed (see subsection 2.2).

In section 3, we prove a structure theorem for $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 (see Theorem 3.9). We further give a full structural description of (P_7, C_4, C_6) -free graphs that contain an induced C_7 and do not admit a clique-cutset; as we shall see, these graphs are precisely the same as $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 and contain no simplicial vertices, which are, in turn, precisely the $(4K_1, C_4, C_6)$ -free graphs that contain an induced C_7 (see Theorem 3.11). The structure of $(4K_1, C_4, C_6)$ -free graphs that contain an induced C_7 was originally described in [14].

In section 4, we give a full structural description of $(2P_3, C_4, C_6, C_7)$ -free graphs that contain an induced T_0 (see Theorem 4.6). As a corollary, we obtain a full structural description of $(2P_3, C_4, C_6)$ -free graphs that contain an induced T_0 and do not contain a simplicial vertex (see Corollary 4.7).

In section 5, we combine the results of sections 3 and 4 to obtain the main structural results of this paper: a full structural description of $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 (see Theorem 5.1), and a full structural description of $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 and contain no simplicial vertices (see Theorem 5.2).

In section 6, we prove that every $(2P_3, C_4, C_6)$ -free graph that contains an induced C_7 or an induced T_0 , either has a simplicial vertex or has clique-width at most 12 (see Theorem 6.6).

Finally, in section 7, we show that $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 can be recognized and colored in $O(n^3)$ time.

2 Preliminaries

2.1 Terminology and notation

When we say that “sets $X_1, \dots, X_\ell$ form a partition of the set X ,” or that “ $(X_1, \dots, X_\ell)$ is a partition of the set X ,” we mean that sets $X_1, \dots, X_\ell$ are pairwise disjoint and that $X = X_1 \cup \dots \cup X_\ell$. However, slightly nonstandardly, we allow some (or all) of the X_i ’s to be empty.

The vertex and edge set of a graph G are denoted by $V(G)$ and $E(G)$, respectively. For a vertex x in a graph G , the *open neighborhood* (or simply *neighborhood*) of x in G , denoted by $N_G(x)$, is the set of all neighbors of x in G ; the *closed neighborhood* of x in G , denoted by $N_G[x]$, is defined as $N_G[x] := \{x\} \cup N_G(x)$; the *degree* of x in G , denoted by $d_G(x)$, is the number of neighbors that x has in G , i.e. $d_G(x) := |N_G(x)|$.

A vertex v of a graph G is *universal* in G if $N_G[v] = V(G)$, i.e. if v is adjacent to all other vertices of the graph G .

Distinct vertices x, y of a graph G are *twins* in G if $N_G[x] = N_G[y]$. (Note that twins are, in particular, adjacent to each other.)

For a graph G and a set $X \subseteq V(G)$, the *open neighborhood* of X in G , denoted by $N_G(X)$, is the set of all vertices in $V(G) \setminus X$ that have a neighbor in X , and the *closed neighborhood* of X in G is the set $N_G[X] := X \cup N_G(X)$.

For a graph G and a nonempty set $X \subseteq V(G)$, we denote by $G[X]$ the subgraph of G induced by X ; for vertices $x_1, \dots, x_t \in V(G)$, we sometimes write $G[x_1, \dots, x_t]$ instead of $G[\{x_1, \dots, x_t\}]$. For a set $X \subsetneq V(G)$, we set $G \setminus X := G[V(G) \setminus X]$. If G has at least two vertices and $x \in V(G)$, we sometimes write $G \setminus x$ instead of $G \setminus \{x\}$.⁸

Given a graph G , a vertex $x \in V(G)$, and a set $Y \subseteq V(G) \setminus \{x\}$, we say that x is *complete* (resp. *anticomplete*) to Y in G provided that x is adjacent (resp. nonadjacent) to all vertices of Y in G , and we say that x is *mixed* on Y in G if x is neither complete nor anticomplete to Y in G (i.e. if x has both a neighbor and a nonneighbor in Y).

Given a graph G and disjoint sets $X, Y \subseteq V(G)$, we say that X is *complete* (resp. *anticomplete*) to Y in G if every vertex in X is complete (resp. anticomplete) to Y in G .

When we say, for graphs G and Q , that “ G can be obtained from Q by possibly adding universal vertices to it,” we mean that Q is an induced subgraph of G , and $V(G) \setminus V(Q)$ is a (possibly empty) clique of G , complete to $V(Q)$ in G . Note that in this case, every vertex in $V(G) \setminus V(Q)$ is a universal vertex of G . However, it is possible that $G = Q$, in which case G contains no universal vertices unless Q does.

A *thickening* of a graph G is a graph G^* for which there exists a family $\{X_v\}_{v \in V(G)}$ of pairwise disjoint, nonempty sets such that all the following hold:

- $V(G^*) = \bigcup_{v \in V(G)} X_v$;
- for all $v \in V(G)$, X_v is a clique of G^* ;
- for all distinct $u, v \in V(G)$, the following hold:
 - if u, v are adjacent in G , then X_u and X_v are complete to each other in G^* ;
 - if u, v are nonadjacent in G , then X_u and X_v are anticomplete to each other in G^* .

When we write, for a graph G , that “ $x_0, \dots, x_t$ is an induced path in G ” ($t \geq 0$), we mean that $x_0, \dots, x_t$ are pairwise distinct vertices, and that the edges of $G[x_0, \dots, x_t]$ are precisely $x_0x_1, x_1x_2, \dots, x_{t-1}x_t$; the *endpoints* of this path are x_0 and x_t (note that if $t = 0$, then our path has only one endpoint), the *interior vertices* of the path are all its vertices other than the endpoints, and furthermore, we say that our path is *between* its endpoints x_0 and x_t . The *length* of a path is the number of edges that it contains; we may also refer to a path of length k as a “ k -edge path” or as a “ $(k+1)$ -vertex path.” A path is *trivial* if it is of length zero (i.e. if it has only one vertex and no edges), and it is *nontrivial* if it is of length at least one.

For an integer $k \geq 4$, a k -hole in a graph G is an induced C_k in G . When we write that “ $x_0, x_1, \dots, x_{k-1}, x_0$ is a k -hole in G ” ($k \geq 4$), we mean that $x_0, x_1, \dots, x_{k-1}$ are pairwise distinct vertices of G , and that the edges of $G[x_0, x_1, \dots, x_{k-1}]$ are precisely the following: $x_0x_1, x_1x_2, \dots, x_{k-2}x_{k-1}, x_{k-1}x_0$.

The *complement* of a graph G , denoted by $\overline{G}$, is the graph whose vertex set is $V(G)$ and in which two distinct vertices are adjacent if and only if they are nonadjacent in G . A graph is *anticonnected* if its complement is connected. Note that every anticonnected graph on at least two vertices contains a pair of nonadjacent vertices.

As usual, a *connected component* (or simply *component*) of a graph G is a maximal connected induced subgraph of G . An *anticomponent* of a graph G is a maximal anticonnected induced

⁸Since our graphs are nonnull, if x is the only vertex of a graph G , then $G \setminus x$ is not defined.

subgraph of G ; clearly, Q is an anticomponent of G if and only if $\overline{Q}$ is a component of $\overline{G}$. An anticomponent is *trivial* if it has only one vertex, and it is *nontrivial* if it has at least two vertices. Clearly, the vertex sets of the anticomponents of a graph G are complete to each other in G . Moreover, note that the unique vertex of any trivial anticomponent of a graph is a universal vertex of that graph.

A *cutset* of a graph G is a (possibly empty) set $C \subsetneq V(G)$ such that $G \setminus C$ is disconnected. A *clique-cutset* of a graph G is a cutset of G that is also a clique of G . (In particular, if G is disconnected, then $\emptyset$ is a clique-cutset of G .) Note that if v is a simplicial vertex of a graph G , then either G is complete, or $N_G(v)$ is a clique-cutset of G .

2.2 A few simple results

Proposition 2.1. *Let G and Q be graphs. Assume that Q is anticonnected and contains at least two vertices. Then the following are equivalent:*

- G contains exactly one nontrivial anticomponent, and that anticomponent is Q ;
- G can be obtained from Q by possibly adding universal vertices to it.

Proof. This follows immediately from the appropriate definitions. $\square$

Proposition 2.2. *Assume that a graph G contains exactly one nontrivial anticomponent, call it Q . Then G contains a simplicial vertex if and only if Q contains a simplicial vertex.*

Proof. Set $W := V(G) \setminus V(Q)$. Since Q is the only nontrivial anticomponent of G , we see that W is a (possibly empty) clique of G , complete to $V(Q)$ in G .

Suppose first that Q contains a simplicial vertex, call it v . Then $N_G(v) = N_Q(v) \cup W$. Since v is simplicial, we know that $N_Q(v)$ is a clique of Q , and therefore of G as well. Since W is a clique, complete to $V(Q)$ in G , we deduce that $N_G(v)$ is a clique of G , and consequently, v is simplicial in G .

Suppose conversely that G contains a simplicial vertex, call it v . Since every vertex of W is complete to $V(Q)$, and since Q contains a pair of distinct, nonadjacent vertices (because Q is a nontrivial anticomponent of G), we see that no vertex of W is simplicial in G . Therefore, $v \notin W$, and we deduce that $v \in V(Q)$. But clearly, $N_Q(v) = N_G(v) \setminus W$; since $N_G(v)$ is a clique of G , it is clear that $N_Q(v)$ is a clique of Q , and we conclude that v is simplicial in Q . $\square$

Proposition 2.3. *Assume that a graph G contains exactly one nontrivial anticomponent, call it Q . Let H be a graph that contains no universal vertices. Then G is H -free if and only if Q is H -free.*

Proof. If G is H -free, then it is clear that Q is also H -free. Suppose now that G is not H -free; we must show that Q is not H -free either. Fix a nonempty set $X \subseteq V(G)$ such that $G[X]$ is isomorphic to H . If $X \subseteq V(Q)$, then $Q[X] = G[X]$, and we deduce that Q is not H -free. We may now assume that $X \not\subseteq V(Q)$. Fix any $x \in X \setminus V(Q)$. Since Q is the only nontrivial anticomponent of G , we know that $V(G) \setminus V(Q)$ is a clique, complete to $V(Q)$ in G , and we deduce that x is a universal vertex of G . But then x is also a universal vertex of $G[X]$, which is impossible since $G[X]$ is isomorphic to H , and H contains no universal vertices. $\square$

Proposition 2.4. *Let G be a C_4 -free graph, let $S \subseteq V(G)$, and let $v_1, v_2 \in V(G) \setminus S$ be distinct, nonadjacent vertices, complete to S in G . Then S is a clique.*

Proof. Suppose otherwise, and fix distinct, nonadjacent vertices $s_1, s_2 \in S$. Then v_1, s_1, v_2, s_2, v_1 is a 4-hole in G , contrary to the fact that G is C_4 -free. $\square$

Proposition 2.5. *A graph is P_3 -free if and only if all its components are complete graphs.*

Proof. Fix a graph G . If all components of G are complete graphs, then it is clear that G is P_3 -free. Suppose, conversely, that G is P_3 -free, and let C be any component of G ; we must show that C is a complete graph. Suppose otherwise, and fix distinct, nonadjacent vertices c, c' of the component C . Since C is connected, we see that it contains an induced path P between c and c' . Since c, c' are nonadjacent, the path P contains at least three vertices. But now any three consecutive vertices of P induce a P_3 in G , a contradiction. $\square$

Proposition 2.6. *Let G be a C_4 -free graph, and let X and Y be disjoint cliques in G . Then all the following hold:*

- (a) *for all $x, x' \in X$, one of $N_G(x) \cap Y$ and $N_G(x') \cap Y$ is a subset of the other;*
- (b) *X can be ordered as $X = \{x_1, \dots, x_{|X|}\}$ so that $N_G(x_{|X|}) \cap Y \subseteq \dots \subseteq N_G(x_1) \cap Y$;*
- (c) *for all $y, y' \in Y$, one of $N_G(y) \cap X$ and $N_G(y') \cap X$ is a subset of the other;*
- (d) *Y can be ordered as $Y = \{y_1, \dots, y_{|Y|}\}$ so that $N_G(y_{|Y|}) \cap X \subseteq \dots \subseteq N_G(y_1) \cap X$.*

Proof. By symmetry, it suffices to prove (a) and (b). Moreover, it is clear that (a) immediately implies (b), and so it is enough to prove (a). Suppose that (a) is false, and fix $x, x' \in X$ such that neither one of $N_G(x) \cap Y$ and $N_G(x') \cap Y$ is a subset of the other. Then there exist $y, y' \in Y$ such that y is adjacent to x and nonadjacent to x' , and such that y' is adjacent to x' and nonadjacent to x . But then x, x', y', y, x is a 4-hole in G , contrary to the fact that G is C_4 -free. $\square$

Proposition 2.7. *Let G be a $(2P_3, C_4)$ -free graph that admits a clique-cutset. Then G contains a simplicial vertex.*

Proof. Fix a clique-cutset C of G , and let A and B be the vertex sets of two distinct components of $G \setminus C$. Then at least one of A, B is a clique, for otherwise, Proposition 2.5 guarantees that $G[A]$ and $G[B]$ both contain an induced P_3 , and consequently (since A and B are anticomplete to each other), G contains an induced $2P_3$, a contradiction. By symmetry, we may assume that A is a clique. Set $t := |A|$; since $A \neq \emptyset$, we know that $t \geq 1$. We now apply Proposition 2.6(b) to the cliques A and C , and we fix an ordering $A = \{a_1, \dots, a_t\}$ of A such that $N_G(a_t) \cap C \subseteq \dots \subseteq N_G(a_1) \cap C$. Since A is a clique, and since $N_G(A) \subseteq C$, we see that $N_G[a_t] = A \cup (N_G(a_t) \cap C)$. But now $N_G[a_t]$ is a clique: indeed, A and C are both cliques, and our ordering of A guarantees that $N_G(a_t) \cap C$ is complete to A . So, a_t is simplicial in G . $\square$

Proposition 2.8. *Let G be a graph, and let $X \subsetneq V(G)$. Assume that for all $x \in X$, the set $V(G) \setminus N_G[x]$ can be partitioned into cliques of G , pairwise anticomplete to each other.⁹ Then G is $2P_3$ -free if and only if $G \setminus X$ is $2P_3$ -free.*

Proof. If G is $2P_3$ -free, then so is its induced subgraph $G \setminus X$. We now assume that G is not $2P_3$ -free, and we show that $G \setminus X$ is not $2P_3$ -free either. Fix pairwise distinct vertices $a_0, a_1, a_2, b_0, b_1, b_2 \in V(G)$ such that a_0, a_1, a_2 and b_0, b_1, b_2 are both induced paths in G , and such that $\{a_0, a_1, a_2\}$ and $\{b_0, b_1, b_2\}$ are anticomplete to each other in G . Clearly, it is enough to show that $X \cap \{a_0, a_1, a_2, b_0, b_1, b_2\} = \emptyset$. Suppose otherwise. By symmetry, we may assume that there exists an index $i \in \{0, 1, 2\}$ such that $a_i \in X$. So, by hypothesis and by Proposition 2.5, we have that $G \setminus N_G[a_i]$ is P_3 -free. But on the other hand, since $\{a_1, a_2, a_3\}$ and $\{b_1, b_2, b_3\}$ are anticomplete to each other, we see that $b_1, b_2, b_3 \in V(G) \setminus N_G[a_i]$, and consequently, $G \setminus N_G[a_i]$ contains an induced P_3 , a contradiction. $\square$

⁹It is possible that $V(G) \setminus N_G[x] = \emptyset$, i.e. that $N_G[x] = V(G)$.

Proposition 2.9. *Let G be a graph, and let $A \subsetneq V(G)$. Assume that A is either empty or can be partitioned into nonempty cliques, pairwise anticomplete to each other. Assume furthermore that A is anticomplete to $V(G) \setminus A$. Then the following hold:*

- (a) *for any graph H such that no component of H is a complete graph, we have that G is H -free if and only if $G \setminus A$ is H -free;*
- (b) *G is $2P_3$ -free if and only if $G \setminus A$ is $2P_3$ -free.*

Proof. Clearly, (a) implies (b). So, it is enough to prove (a). Fix a graph H such that no component of H is a complete graph. If G is H -free, then $G \setminus A$ is obviously H -free as well. Now, suppose that G is not H -free; we must show that $G \setminus A$ is not H -free either. Fix $X \subseteq V(G)$ such that $G[X]$ is isomorphic to H . It now suffices to show that $X \cap A = \emptyset$, for this will immediately imply that $G \setminus A$ is not H -free, which is what we need to show. Suppose otherwise, that is, suppose that $X \cap A \neq \emptyset$. This, in particular, implies that $A \neq \emptyset$, and so by hypothesis, A can be partitioned into nonempty cliques $A_1, \dots, A_\ell$, pairwise anticomplete to each other, and anticomplete to $V(G) \setminus A$. By symmetry, we may assume that $X \cap A_1 \neq \emptyset$. Since A_1 is a clique, anticomplete to $V(G) \setminus A_1$, we deduce that $G[X \cap A_1]$ is a complete graph that is also a component of $G[X]$. But this is impossible since $G[X]$ is isomorphic to H , and no component of H is a complete graph. $\square$

Proposition 2.10. *Let G^* be a thickening of a graph G . Then all the following hold:*

- (a) *G^* contains a simplicial vertex if and only if G contains a simplicial vertex;*
- (b) *G^* contains a universal vertex if and only if G contains a universal vertex;*
- (c) *if G has at least two vertices, then G^* is anticonnected if and only if G is anticonnected;*
- (d) *for any graph H that contains no pair of twins, G^* is H -free if and only if G is H -free.*

Proof. Using the fact that G^* is a thickening of G , we fix a family $\{X_v\}_{v \in V(G)}$ of pairwise disjoint, nonempty sets such that all the following hold:

- $V(G^*) = \bigcup_{v \in V(G)} X_v$;
- for all $v \in V(G)$, X_v is a clique of G^* ;
- for all distinct $u, v \in V(G)$, the following hold:
 - if u, v are adjacent in G , then X_u and X_v are complete to each other in G^* ;
 - if u, v are nonadjacent in G , then X_u and X_v are anticomplete to each other in G^* .

We first prove (a-b). If v is a simplicial (respectively: universal) vertex of G , then clearly, any vertex of X_v is simplicial (respectively: universal) in G^* . On the other hand, if a vertex x is simplicial (respectively: universal) of G^* , then we fix $v \in V(G)$ such that $x \in X_v$, and we observe that v is a simplicial (respectively: universal) vertex of G . This proves (a-b).

Next, we prove (c). Assume that G has at least two vertices. We will prove that G^* is not anticonnected if and only if G is not anticonnected.

Suppose first that G is not anticonnected. Then $V(G)$ can be partitioned into nonempty sets A and B that are complete to each other in G . But then $\bigcup_{v \in A} X_v$ and $\bigcup_{v \in B} X_v$ form a partition of $V(G^*)$ into nonempty sets, complete to each other in G^* , and we deduce that G^* is not anticonnected.

Suppose now that G^* is not anticonnected. Then $V(G^*)$ can be partitioned into nonempty sets A^* and B^* , complete to each other in G^* . If there exists some $v \in V(G)$ such that X_v intersects both A^* and B^* , then v is a universal vertex of G ,¹⁰ which (since G has at least two vertices) implies that G is not anticonnected. So, assume that no such vertex v exists. Then $A := \{v \in V(G) \mid X_v \subseteq A^*\}$ and $B := \{v \in V(G) \mid X_v \subseteq B^*\}$ form a partition of $V(G)$ into nonempty sets, complete to each other in G . So, G is not anticonnected. This proves (c).

It remains to prove (d). Fix a graph H that contains no pair of twins. Clearly, G is (isomorphic to) an induced subgraph of G^* , and so if G^* is H -free, so is G . Now, suppose that G^* is not H -free; we must show that G is not H -free either. Fix $X \subseteq V(G^*)$ such that $G^*[X]$ is isomorphic to H . Note that for all $v \in V(G)$, any two distinct vertices of X_v are twins in G^* . Since H contains no pair of twins, we deduce that $|X \cap X_v| \leq 1$ for all $v \in V(G)$. Now, for all $v \in V(G)$ such that $|X \cap X_v| = 1$, let v^* be the unique vertex of $X \cap X_v$. On the other hand, for all $v \in V(G)$ such that $X \cap X_v = \emptyset$, let v^* be an arbitrarily chosen vertex of X_v . Set $X^* := \{v^* \mid v \in V(G)\}$. Then $G[X^*]$ is isomorphic to G and contains $G^*[X]$ as an induced subgraph. Since $G^*[X]$ is isomorphic to H , we deduce that G is not H -free. This proves (d). $\square$

3 (P_7, C_4, C_6) -free graphs that contain an induced C_7

In this section, we give a full structural description of (P_7, C_4, C_6) -free graphs that contain an induced C_7 and do not admit a clique-cutset (see Theorem 3.11). As mentioned in the introduction, it turns out that these graphs are exactly the same as the $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 and contain no simplicial vertices, which are in turn exactly the same as the $(4K_1, C_4, C_6)$ -free graphs that contain an induced C_7 . We remind the reader that the structure of $(4K_1, C_4, C_6)$ -free graphs that contain an induced C_7 was originally described in [14].

In this section, we further fully describe the structure of all $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 (see Theorem 3.9, and note that this theorem also applies to graphs that may possibly contain simplicial vertices).

3.1 The special partition and the 7-saucer

A *special partition* of a graph G is a partition $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ of $V(G)$ into cliques such that all the following are satisfied:¹¹

- (a) cliques $X_0, \dots, X_6$ are all nonempty;¹²
- (b) for all $i \in \mathbb{Z}_7$, X_i is complete to X_{i-1}, X_{i+1} and anticomplete to $X_{i+2}, X_{i+3}, X_{i+4}, X_{i+5}$;
- (c) for all $i \in \mathbb{Z}_7$, X_i is complete to $Y_i, Y_{i+3}, Y_{i+6}, Z_i, Z_{i+3}, Z_{i+4}, Z_{i+5}, Z_{i+6}, W$ and anticomplete to $Y_{i+1}, Y_{i+2}, Y_{i+4}, Y_{i+5}, Z_{i+1}, Z_{i+2}$;
- (d) for all $i \in \mathbb{Z}_7$, if $Y_i \neq \emptyset$, then $Y_{i+1}, Y_{i+2}, Y_{i+5}, Y_{i+6}, Z_{i+5}, Z_{i+6}$ are all empty, and at most one of Y_{i+3}, Y_{i+4} is nonempty;¹³

¹⁰Indeed, fix any $u \in V(G) \setminus \{v\}$. Then X_u intersects at least one of A^* and B^* . By symmetry, we may assume that $X_u \cap A^* \neq \emptyset$. Since A^* is complete to B^* in G^* , we know that $X_u \cap A^*$ is complete to $X_v \cap B^*$ in G^* . Since $X_u \cap A^*$ and $X_v \cap B^*$ are both nonempty, we see that there is at least one edge between X_u and X_v in G^* . So, by definition, u is adjacent to v .

¹¹Indices are understood to be in $\mathbb{Z}_7$.

¹²Cliques $Y_0, \dots, Y_6, Z_0, \dots, Z_6, W$ may possibly be empty.

¹³In particular, at most two of the cliques $Y_0, \dots, Y_6$ are nonempty.

- (e) for all $i \in \mathbb{Z}_7$, if $Z_i \neq \emptyset$, then Z_{i+2}, Z_{i+5} are empty;¹⁴
- (f) for all $i \in \mathbb{Z}_7$, Y_i is complete to $Y_{i+3}, Y_{i+4}, Z_i, Z_{i+1}, Z_{i+3}, Z_{i+4}, W$ and anticomplete to Z_{i+2} ;
- (g) for all $i \in \mathbb{Z}_7$, Z_i is complete to $Z_{i+1}, Z_{i+3}, Z_{i+4}, Z_{i+6}, W$.

Note that if $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of a graph G , then by definition, the following hold:

- for all $i \in \mathbb{Z}_7$, Y_i is complete to X_i, X_{i+1}, X_{i+4} and anticomplete to $X_{i+2}, X_{i+3}, X_{i+5}, X_{i+6}$;
- for all $i \in \mathbb{Z}_7$, Z_i is complete to $X_i, X_{i+1}, X_{i+2}, X_{i+3}, X_{i+4}$ and anticomplete to X_{i+5}, X_{i+6} .
- the clique W is complete to $V(G) \setminus W = (X_0 \cup \dots \cup X_6) \cup (Y_0 \cup \dots \cup Y_6) \cup (Z_0 \cup \dots \cup Z_6)$, and consequently, all vertices in W are universal vertices of G .¹⁵

The following lemma may help the reader gain a more intuitive understanding of the special partition. (For further intuition regarding the special partition, see the definition of the family $\mathcal{M}$ in subsection 3.2, and then see Proposition 3.6.)

Lemma 3.1 (Lemma 7.2 of [23]). *Let $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ be a special partition of a graph G , and set $Y := \bigcup_{i \in \mathbb{Z}_7} Y_i$ and $Z := \bigcup_{i \in \mathbb{Z}_7} Z_i$. Then exactly one of the following holds:*

- (a) *there exists an index $i \in \mathbb{Z}_7$ such that $Y = Y_i \cup Y_{i+3}$ and $Z = Z_i \cup Z_{i+3} \cup Z_{i+4}$;*
- (b) *there exists an index $i \in \mathbb{Z}_7$ such that all the following hold:*

- $Y = Y_i$ and $Z = Z_{i+1} \cup Z_{i+2} \cup Z_{i+3}$,
- Y_i, Z_{i+2} are both nonempty,
- at most one of Z_{i+1}, Z_{i+3} is nonempty.

Moreover, all the following hold:

- Y and Z are both cliques;
- if (a) holds, then $Y \cup Z$ is a clique;
- if (b) holds, then $Y \cup Z$ is not a clique.

It was shown in [14] that every $(4K_1, C_4, C_6)$ -graph that contains an induced C_7 admits a special partition.¹⁶ However, our proof does not actually rely on this result. Instead, we only use the converse, obtained via routine case checking in [23] and stated below.

Lemma 3.2 (Lemma 7.3(a) of [23]). *Any graph that admits a special partition is $(4K_1, C_4, C_6)$ -free and contains an induced C_7 .*

A *7-saucer* is a graph G whose vertex set can be partitioned into sets $X_0, \dots, X_6, Y_0, \dots, Y_6, Z_0, \dots, Z_6, W, A$ (with indices understood to be in $\mathbb{Z}_7$) such that all the following hold:

- $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of $G \setminus A$;

¹⁴In particular, at most three of the cliques $Z_0, \dots, Z_6$ are nonempty.

¹⁵However, it is possible that $W = \emptyset$, in which case, G contains no universal vertices.

¹⁶The term “special partition” was actually introduced in [23], not in [14]. However, as the reader can check, the structure described in [14] is precisely our “special partition.”

- A is anticomplete to $X_0, \dots, X_6$;¹⁷
- for all $i \in \mathbb{Z}_7$, either A is anticomplete to Y_i , or $Z_{i+2} = \emptyset$;
- either $A = \emptyset$, or A can be partitioned into nonempty cliques $A_1, \dots, A_\ell$, pairwise anticomplete to each other, such that for all $i \in \{1, \dots, \ell\}$, A_i can be ordered as $A_i = \{a_1^i, \dots, a_{r_i}^i\}$ so that $N_G[a_{r_i}^i] \subseteq \dots \subseteq N_G[a_1^i]$.

Under these circumstances, we also say that $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$ is a *7-saucer partition* of the 7-saucer G .

One of the main goals of this section is to prove that a graph is $(2P_3, C_4, C_6)$ -free and contains an induced C_7 if and only if it is a 7-saucer (see Theorem 3.9; for the “if” part, see Proposition 3.4, and for the “only if” part, see Lemma 3.8).

Proposition 3.3. *Let G be a 7-saucer, and let $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$ be an associated 7-saucer partition of G . Set $X := \bigcup_{i \in \mathbb{Z}_7} X_i$, $Y := \bigcup_{i \in \mathbb{Z}_7} Y_i$ and $Z := \bigcup_{i \in \mathbb{Z}_7} Z_i$. Then all the following hold:*

- (a) $N_G(A) \subseteq Y \cup Z \cup W$;
- (b) $N_G(A)$ is a clique;
- (c) for all $v \in N_G(A)$, the set $V(G) \setminus N_G[v]$ can be partitioned into cliques, pairwise anticomplete to each other.¹⁸

Proof. In our proof, we will repeatedly use the fact that $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of $G \setminus A$, as per the definition of a 7-saucer.

For (a), we simply observe that A is anticomplete to X (by the definition of a 7-saucer), and so, $N_G(A) \subseteq Y \cup Z \cup W$.

Next, we prove (b). If $Y \cup Z \cup W$ is a clique, then we are done by (a). We may therefore assume that $Y \cup Z \cup W$ is not a clique. By the definition of a special partition, W is a clique, complete to $X \cup Y \cup Z$. So, since $Y \cup Z \cup W$ is not a clique, we know that $Y \cup Z$ is not a clique. By Lemma 3.1 and by symmetry, we may now assume that all the following hold:

- $Y = Y_0$ and $Z = Z_1 \cup Z_2 \cup Z_3$,
- Y_0, Z_2 are both nonempty,
- at most one of Z_1, Z_3 is nonempty.

Since Z_2 is nonempty, the definition of a 7-saucer guarantees that A is anticomplete to Y_0 . We now deduce that $N_G(A) \subseteq Z_1 \cup Z_2 \cup Z_3 \cup W$. But by the definition of a special partition, $Z_1 \cup Z_2 \cup Z_3 \cup W$ is a clique; consequently, $N_G(A)$ is also a clique. This proves (b).

It remains to prove (c). Fix a vertex $v \in N_G(A)$. We must show that $V(G) \setminus N_G[v]$ can be partitioned into cliques, pairwise anticomplete to each other. Recall that $N_G(A) \subseteq Y \cup Z \cup W$; thus, $v \in Y \cup Z \cup W$.

Suppose first that $v \in W$. By the definitions of a 7-saucer and a special partition, W is a clique, complete to $V(G) \setminus (W \cup A)$. So, $V(G) \setminus N_G[v] \subseteq A$. But by the definition of a 7-saucer, A can be partitioned into (possibly empty) cliques, pairwise anticomplete to each other.

Next, suppose that $v \in Y$. By symmetry, we may assume that $v \in Y_0$. Since $v \in N_G(A) \cap Y_0$, we see that A is not anticomplete to Y_0 , and so by the definition of a 7-saucer, we know that $Z_2 =$

¹⁷Thus, $N_G(A) \subseteq (Y_0 \cup \dots \cup Y_6) \cup (Z_0 \cup \dots \cup Z_6) \cup W$.

¹⁸Here, it is possible that $V(G) \setminus N_G[v] = \emptyset$, i.e. that $V(G) = N_G[v]$, so that v is a universal vertex of G .

$\emptyset$. Then the definitions of a 7-saucer and a special partition imply that Y_0 is a clique, complete to $(X_0 \cup X_1 \cup X_4)(Y \setminus Y_0) \cup Z \cup W$, and furthermore, that $V(G) \setminus N_G[v] \subseteq (X_2 \cup X_3) \cup (X_5 \cup X_6) \cup A$. But once again by the definitions of a 7-saucer and a special partition, all the following hold:

- A can be partitioned into (possibly empty) cliques, pairwise anticomplete to each other;
- A is anticomplete to X (and therefore, to $X_2 \cup X_3$ and $X_5 \cup X_6$ as well);
- $X_2 \cup X_3$ and $X_5 \cup X_6$ are disjoint cliques, anticomplete to each other.

Thus, $V(G) \setminus N_G[v]$ can be partitioned into cliques, pairwise anticomplete to each other.

Finally, suppose that $v \in Z$. By symmetry, we may assume that $v \in Z_2$. In particular, $Z_2 \neq \emptyset$, and so by the definition of a 7-saucer, A is anticomplete to Y_0 . On the other hand, the definition of a special partition implies that Z_2 is a clique, complete to $(X_2 \cup X_3 \cup X_4 \cup X_5 \cup X_6) \cup (Y \setminus Y_0) \cup (Z \setminus Z_2) \cup W$. Thus, $V(G) \setminus N_G[v] \subseteq X_0 \cup X_1 \cup Y_0 \cup A$. It is now enough to show that $X_0 \cup X_1 \cup Y_0 \cup A$ can be partitioned into cliques, pairwise anticomplete to each other. By the definition of a special partition, $X_0 \cup X_1 \cup Y_0$ is a clique. Next, the definition of a 7-saucer guarantees that A is anticomplete to $X_0 \cup X_1$, and we already saw that A is anticomplete to Y_0 . Finally, the definition of a special partition guarantees that A can be partitioned into (possibly empty) cliques, pairwise anticomplete to each other. This completes the proof of (c). $\square$

Proposition 3.4. *Every 7-saucer is $(2P_3, C_4, C_6)$ -free and contains an induced C_7 .*

Proof. Fix a 7-saucer G , and let $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$ be an associated 7-saucer partition of G . By definition, $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of $G \setminus A$, and so by Lemma 3.2, $G \setminus A$ is $(4K_1, C_4, C_6)$ -free and contains an induced C_7 . Since $G \setminus A$ contains an induced C_7 , so does G . It remains to show that G is $(2P_3, C_4, C_6)$ -free.

Claim 3.4.1. *Every hole of G is in fact a hole of $G \setminus A$. Consequently, G is (C_4, C_6) -free.*

Proof of Claim 3.4.1. Fix a hole $v_0, v_1, \dots, v_{k-1}, v_0$ (with $k \geq 4$, and with indices understood to be in $\mathbb{Z}_k$) in G . We must show that $\{v_0, v_1, \dots, v_{k-1}\} \cap A = \emptyset$. Suppose otherwise. Let B be the vertex set of a component of $G[A]$ such that $\{v_0, v_1, \dots, v_{k-1}\} \cap B \neq \emptyset$. By the definition of a 7-saucer, B is a clique, and there exists an ordering $B = \{b_1, \dots, b_r\}$ of B such that $N_G[b_r] \subseteq \dots \subseteq N_G[b_1]$. Since $N_G(A)$ is clique (by Proposition 3.3), we see that for all $i \in \{1, \dots, r\}$, b_i is simplicial in $G \setminus \{b_{i+1}, \dots, b_r\}$. Now, fix the largest index $i \in \{1, \dots, r\}$ such that $b_i \in \{v_0, v_1, \dots, v_{k-1}\}$; by symmetry, we may assume that $b_i = v_0$. Then $v_0, v_1, \dots, v_{k-1}, v_0$ is a hole in $G \setminus \{b_{i+1}, \dots, b_r\}$, and we deduce that $v_0 = b_i$ is a simplicial vertex in the hole $v_0, v_1, \dots, v_{k-1}, v_0$, contrary to the fact that holes contain no simplicial vertices. This proves that every hole of G is in fact a hole of $G \setminus A$. Since $G \setminus A$ is (C_4, C_6) -free, we deduce that G is also (C_4, C_6) -free. $\blacklozenge$

Claim 3.4.2. *G is $2P_3$ -free.*

Proof of Claim 3.4.2. By Propositions 2.8 and 3.3, it suffices to show that $G' := G \setminus N_G(A)$ is $2P_3$ -free. First, by the definition of a 7-saucer, A is either empty or can be partitioned into cliques, pairwise anticomplete to each other. Moreover, it is clear that A is anticomplete to $V(G') \setminus A = V(G) \setminus N_G[A]$ in G' . So, by Proposition 2.9(b), it is enough to show that $G' \setminus A = G \setminus N_G[A]$ is $2P_3$ -free; since $G' \setminus A = G \setminus N_G[A]$ is an induced subgraph of $G \setminus A$, and since $4K_1$ is an induced subgraph of $2P_3$, it in fact suffices to show that $G \setminus A$ is $4K_1$ -free. But this follows from Lemma 3.2, since by the definition of a 7-saucer, $G \setminus A$ admits a special partition, namely $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$. $\blacklozenge$

Claims 3.4.1 and 3.4.2 together guarantee that G is $(2P_3, C_4, C_6)$ -free, which completes the argument. $\square$

3.2 The family $\mathcal{M}$

In this subsection, we define the family $\mathcal{M}$ of graphs, which we need in order to properly state Theorem 3.11, our structure theorem for (P_7, C_4, C_6) -free graphs that contain an induced C_7 and do not admit a clique-cutset. We first define graphs M_0 , M_1 , M_2 , and M_3 , as follows.

M_0 is the graph with vertex set $\{x_0, x_1, x_2, x_3, x_4, x_5, x_6, y_0, y_3, z_0, z_3, z_4\}$ and with adjacency as follows:

- $x_0, x_1, x_2, x_3, x_4, x_5, x_6$ is a 7-hole;
- y_0 is complete to $\{x_0, x_1, x_4\}$ and anticomplete to $\{x_2, x_3, x_5, x_6\}$;
- y_3 is complete to $\{x_3, x_4, x_0\}$ and anticomplete to $\{x_1, x_2, x_5, x_6\}$;
- z_0 is complete to $\{x_0, x_1, x_2, x_3, x_4\}$ and anticomplete to $\{x_5, x_6\}$;
- z_3 is complete to $\{x_3, x_4, x_5, x_6, x_0\}$ and anticomplete to $\{x_1, x_2\}$;
- z_4 is complete to $\{x_4, x_5, x_6, x_0, x_1\}$ and anticomplete to $\{x_2, x_3\}$;
- $\{y_0, y_3, z_0, z_3, z_4\}$ is a clique.

M_1 is the graph with vertex set $\{x_0, x_1, x_2, x_3, x_4, x_5, x_6, y_0, z_2\}$ and with adjacency as follows:

- $x_0, x_1, x_2, x_3, x_4, x_5, x_6$ is a 7-hole;
- y_0 is complete to $\{x_0, x_1, x_4\}$ and anticomplete to $\{x_2, x_3, x_5, x_6\}$;
- z_2 is complete to $\{x_2, x_3, x_4, x_5, x_6\}$ and anticomplete to $\{x_0, x_1\}$;
- y_0 and z_2 are nonadjacent.

M_2 is the graph with vertex set $\{x_0, x_1, x_2, x_3, x_4, x_5, x_6, y_0, z_1, z_2\}$ and with adjacency as follows:

- $M_2 \setminus z_1 = M_1$;
- z_1 is complete to $\{x_1, x_2, x_3, x_4, x_5\}$ and anticomplete to $\{x_6, x_0\}$;
- z_1 is complete to $\{y_0, z_2\}$.

M_3 is the graph with vertex set $\{x_0, x_1, x_2, x_3, x_4, x_5, x_6, y_0, z_2, z_3\}$ and with adjacency as follows:

- $M_3 \setminus z_3 = M_1$;
- z_3 is complete to $\{x_3, x_4, x_5, x_6, x_0\}$ and anticomplete to $\{x_1, x_2\}$;
- z_3 is complete to $\{y_0, z_2\}$.

We now define the family $\mathcal{M}$, as follows:

$$\mathcal{M} := \{M_0 \setminus S \mid S \subseteq \{y_0, y_3, z_0, z_3, z_4\}\} \cup \{M_1, M_2, M_3\}.$$

Thus, $\mathcal{M}$ is the collection of graphs whose members are M_1, M_2, M_3 , plus all induced subgraphs of M_0 that still contain the hole $x_0, x_1, x_2, x_3, x_4, x_5, x_6, x_0$. We note that each graph in $\mathcal{M}$ has at most 12 vertices and contains an induced C_7 .

Proposition 3.5. *Every graph in $\mathcal{M}$ is anticonnected and contains no simplicial vertices and no universal vertices. Consequently, any thickening of a graph in $\mathcal{M}$ is anticonnected contains no simplicial vertices and no universal vertices.*

Proof. In view of Proposition 2.10, it is enough to prove the first statement. So, fix any graph $M \in \mathcal{M}$; we must show that M is anticonnected and contains no simplicial vertices and no universal vertices. First of all, $x_0, x_1, \dots, x_6, x_0$ is a 7-hole in M . Since $\overline{C_7}$ is connected, we see that $M[x_0, x_1, \dots, x_6]$ is anticonnected. Moreover, by the definition of $\mathcal{M}$, every vertex in $V(M) \setminus \{x_0, x_1, \dots, x_6\}$ has a nonneighbor in $\{x_0, x_1, \dots, x_6\}$. Thus, M is anticonnected. Since M has more than one vertex, this implies, in particular, that M contains no universal vertices. It remains to show that M contains no simplicial vertices. First, for all $i \in \mathbb{Z}_7$, x_i is complete to $\{x_{i-1}, x_{i+1}\}$, and so since x_{i-1}, x_{i+1} are nonadjacent, x_i is not a simplicial vertex of M . On the other hand, by the definition of the family $\mathcal{M}$, every vertex in $V(M) \setminus \{x_0, x_1, \dots, x_6\}$ has at least three neighbors in $\{x_0, x_1, \dots, x_6\}$; since $M[x_0, x_1, \dots, x_6]$ is triangle-free, we deduce that no vertex in $V(M) \setminus \{x_0, x_1, \dots, x_6\}$ is simplicial. Therefore, M contains no simplicial vertices. $\square$

Proposition 3.6. *For any graph G , the following are equivalent:*

- (a) G admits a special partition;
- (b) G has exactly one nontrivial anticomponent, and this anticomponent is a thickening of a graph in $\mathcal{M}$;
- (c) G can be obtained from a thickening of a graph in $\mathcal{M}$ by possibly adding universal vertices to it.

Proof. By Proposition 3.5, any thickening of a graph in $\mathcal{M}$ is anticonnected. So, by Proposition 2.1, (b) and (c) are equivalent. It remains to prove that (a) and (c) are equivalent.

Suppose first that (a) holds. Fix a special partition $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ of G . Then W is a clique, complete to $V(G) \setminus W$. Thus, G can be obtained from $G \setminus W$ by possibly adding universal vertices to it. But Lemma 3.1 and the relevant definitions immediately imply that $G \setminus W$ is a thickening of a graph in $\mathcal{M}$. So, (c) holds.

Suppose conversely that (c) holds, so that G can be obtained from a thickening M^* of a graph $M \in \mathcal{M}$ by possibly adding universal vertices to M^* . By the definition of $\mathcal{M}$ and of the special partition, it is clear that M^* admits a special partition $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; \emptyset)$. But now for $W := V(G) \setminus V(M^*)$, we have that W is a clique, complete to $V(M^*) = V(G) \setminus W$, and we deduce that $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of G . So, (a) holds. $\square$

3.3 Decomposing (P_7, C_4, C_6) -free graphs that contain an induced C_7

Lemma 3.7. *Let G be a (P_7, C_4, C_6) -free graph, let $x_0, x_1, x_2, \dots, x_6, x_0$ (with indices in $\mathbb{Z}_7$) be a 7-hole in G , and let $v \in V(G) \setminus \{x_0, x_1, x_2, \dots, x_6\}$. Then one of the following holds:*

- (a) v is anticomplete to $\{x_0, x_1, x_2, \dots, x_6\}$;
- (b) there exists an index $i \in \mathbb{Z}_7$ such that $N_G(v) \cap \{x_0, x_1, x_2, \dots, x_6\} = \{x_{i-1}, x_i, x_{i+1}\}$;
- (c) there exists an index $i \in \mathbb{Z}_7$ such that $N_G(v) \cap \{x_0, x_1, x_2, \dots, x_6\} = \{x_i, x_{i+1}, x_{i+4}\}$;
- (d) there exists an index $i \in \mathbb{Z}_7$ such that $N_G(v) \cap \{x_0, x_1, x_2, \dots, x_6\} = \{x_i, x_{i+1}, x_{i+2}, x_{i+3}, x_{i+4}\}$;
- (e) v is complete to $\{x_0, x_1, x_2, \dots, x_6\}$.

Proof. To simplify notation, set $X := \{x_0, x_1, x_2, \dots, x_6\}$. If v is anticomplete to X , then (a) holds, and we are done. Next, if v has exactly one neighbor in X , then by symmetry, we may assume that $N_G(v) \cap X = \{x_0\}$, and we see that $v, x_0, x_1, x_2, x_3, x_4, x_5$ is an induced 7-vertex path in G , contrary to the fact that G is P_7 -free. Next, suppose that v has exactly two neighbors in X ; by symmetry, we may assume that $N_G(v) \cap X = \{x_0, x_k\}$ for some $k \in \{1, 2, 3\}$. But if $k = 1$, then $v, x_1, x_2, x_3, x_4, x_5, x_6$ is an induced 7-vertex path in G ; if $k = 2$, then v, x_0, x_1, x_2, v is a 4-hole in G ; and if $k = 3$, then $v, x_3, x_4, x_5, x_6, x_0, v$ is a 6-hole in G . Since G is (P_7, C_4, C_6) -free, none of these three outcomes is possible.

From now on, we may assume that $3 \leq |N_G(v)| \leq 6$. Let us define a *sector* to be a nontrivial induced path in the hole $x_0, x_1, \dots, x_6, x_0$ such that v is adjacent to both endpoints of the path and is anticomplete to the interior of the path. The *length* of a sector is simply its length as a path (i.e. the number of edges that it contains). Note that v , together with the vertices of any sector of length ℓ , induces a cycle of length $\ell + 2$ in G . Since G is (P_7, C_4, C_6) -free, we know that all induced cycles of G are of length 3, 5, or 7; therefore, all sectors are of length 1, 3, or 5. Meanwhile, since every edge of the hole $x_0, x_1, \dots, x_6, x_0$ belongs to exactly one sector, we see that the sum of lengths of all the sectors is 7. So, if the number of sectors of length 1, 3, and 5 is a , b , and c , respectively, then we have that $a + 3b + 5c = 7$. But the equation $a + 3b + 5c = 7$ has (exactly) the following four nonnegative integer solutions:

- (1) $a = 2, b = 0, c = 1$;
- (2) $a = 1, b = 2, c = 0$;
- (3) $a = 4, b = 1, c = 0$;
- (4) $a = 7, b = 0, c = 0$.

Clearly, these four solutions correspond (respectively) to the outcomes (b), (c), (d), (e) of the lemma. This completes the argument. $\square$

Lemma 3.8. *Let G be a (P_7, C_4, C_6) -free graph that contains an induced C_7 . Then either G admits a special partition, or G admits a clique-cutset. Moreover, if G is additionally $2P_3$ -free, then G is a 7-saucer.*

Proof. Let $x_0, x_1, x_2, x_3, x_4, x_5, x_6, x_0$ (with indices in $\mathbb{Z}_7$) be a 7-hole in G , and set $\widehat{X} := \{x_0, x_1, x_2, \dots, x_6\}$. Let W be the set of all vertices in $V(G) \setminus \widehat{X}$ that are complete to $\widehat{X}$, let A be the set of all vertices in $V(G) \setminus \widehat{X}$ that are anticomplete to $\widehat{X}$, and for all $i \in \mathbb{Z}_7$, set

- $X_i := \{x_i\} \cup \{v \in V(G) \setminus \widehat{X} \mid N_G(v) \cap \widehat{X} = \{x_{i-1}, x_i, x_{i+1}\}\}$;
- $Y_i := \{v \in V(G) \setminus \widehat{X} \mid N_G(v) \cap \widehat{X} = \{x_i, x_{i+1}, x_{i+4}\}\}$;
- $Z_i := \{v \in V(G) \setminus \widehat{X} \mid N_G(v) \cap \widehat{X} = \{x_i, x_{i+1}, x_{i+2}, x_{i+3}, x_{i+4}\}\}$.

Furthermore, set $X := \bigcup_{i \in \mathbb{Z}_7} X_i$, $Y := \bigcup_{i \in \mathbb{Z}_7} Y_i$, and $Z := \bigcup_{i \in \mathbb{Z}_7} Z_i$.

By Lemma 3.7, $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$ is a partition of $V(G)$.¹⁹ Our goal is to prove all the following:

- $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of $G \setminus A$ (see Claim 3.8.7);
- if $A \neq \emptyset$, then G admits a clique-cutset (see Claim 3.8.9);
- if G is $2P_3$ -free, then G is a 7-saucer (see Claim 3.8.10).

Our first goal is to show that $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of $G \setminus A$. The fact that $X_0, \dots, X_6$ are all nonempty follows from the construction. Further, all the following hold:

- for all $i \in \mathbb{Z}_7$, distinct, nonadjacent vertices x_{i-1}, x_{i+1} are complete to X_i ;
- for all $i \in \mathbb{Z}_7$, distinct, nonadjacent vertices x_i, x_{i+4} are complete to Y_i ;
- for all $i \in \mathbb{Z}_7$, distinct, nonadjacent vertices x_i, x_{i+2} are complete to Z_i ;
- distinct, nonadjacent vertices x_0, x_2 are complete to W .

In view of Proposition 2.4, this implies that $X_0, \dots, X_6, Y_0, \dots, Y_6, Z_0, \dots, Z_6, W$ are all cliques. It remains to check that items (b)-(g) from the definition of a special partition are satisfied. We prove this via a sequence of claims (see Claims 3.8.1-3.8.6 below).

Claim 3.8.1. *For all $i \in \mathbb{Z}_7$, X_i is complete to X_{i-1}, X_{i+1} and anticomplete to $X_{i+2}, X_{i+3}, X_{i+4}, X_{i+5}$.*

Proof of Claim 3.8.1. By symmetry, it suffices show that X_0 is complete to X_1 and anticomplete to X_2, X_3 .

If some $x'_0 \in X_0$ and $x'_1 \in X_1$ are nonadjacent, then $x'_1, x_2, x_3, x_4, x_5, x_6, x'_0$ is an induced 7-vertex path in G , contrary to the fact that G is P_7 -free. So, X_0 is complete to X_1 .

If some $x'_0 \in X_0$ is adjacent to some $x'_2 \in X_2$, then $x'_0, x'_2, x_3, x_4, x_5, x_6, x'_0$ is a 6-hole in G , contrary to the fact that G is C_6 -free. Thus, X_0 is anticomplete to X_2 .

If some $x'_0 \in X_0$ is adjacent to some $x'_3 \in X_3$, then $x'_0, x_1, x_2, x'_3, x'_0$ is a 4-hole in G , contrary to the fact that G is C_4 -free. Thus, X_0 is anticomplete to X_3 . ♦

Claim 3.8.2. *For all $i \in \mathbb{Z}_7$, X_i is complete to $Y_i, Y_{i+3}, Y_{i+6}, Z_i, Z_{i+3}, Z_{i+4}, Z_{i+5}, Z_{i+6}, W$ and anticomplete to $Y_{i+1}, Y_{i+2}, Y_{i+4}, Y_{i+5}, Z_{i+1}, Z_{i+2}$.*

Proof of Claim 3.8.2. By symmetry, it suffices to prove the claim for $i = 0$. If some x'_0 has a nonneighbor $v \in Y_0 \cup Y_3 \cup Y_6 \cup Z_0 \cup Z_3 \cup Z_4 \cup Z_5 \cup Z_6 \cup W$, then v has either exactly two, or exactly four, or exactly six neighbors in the 7-hole $x'_0, x_1, x_2, x_3, x_4, x_5, x_6, x'_0$, contrary to Lemma 3.7. Thus, X_0 is complete to $Y_0, Y_3, Y_6, Z_0, Z_3, Z_4, Z_5, Z_6, W$.

Similarly, if some $x'_0 \in X_0$ has a neighbor $v \in Y_1 \cup Y_2 \cup Y_4 \cup Y_5 \cup Z_1 \cup Z_2$, then v has either exactly four or exactly six neighbors in the 7-hole $x'_0, x_1, x_2, x_3, x_4, x_5, x_6, x'_0$, contrary to Lemma 3.7. Thus, X_0 is anticomplete to $Y_1, Y_2, Y_4, Y_5, Z_1, Z_2$. ♦

Claim 3.8.3. *For all $i \in \mathbb{Z}_7$, if $Y_i \neq \emptyset$, then $Y_{i+1}, Y_{i+2}, Y_{i+5}, Y_{i+6}, Z_{i+5}, Z_{i+6}$ are all empty, and at most one of Y_{i+3}, Y_{i+4} is nonempty.*

¹⁹Consequently, sets X, Y, Z, W, A also form a partition of $V(G)$.

Proof of Claim 3.8.3. By symmetry, it suffices to prove the claim for $i = 0$. So, assume that $Y_0 \neq \emptyset$. We must show that $Y_1, Y_2, Y_5, Y_6, Z_5, Z_6$ are all empty, and that at most one of Y_3, Y_4 is nonempty. Fix $y_0 \in Y_0$.

Suppose that $Y_1 \neq \emptyset$, and fix $y_1 \in Y_1$. If y_0 is adjacent to y_1 , then y_0, y_1, x_5, x_4, y_0 is a 4-hole in G , contrary to the fact that G is C_4 -free. On the other hand, if y_0 is nonadjacent to y_1 , then $y_1, x_2, x_3, x_4, y_0, x_0, x_6$ is an induced 7-vertex path in G , contrary to the fact that G is P_7 -free. This proves that $Y_1 = \emptyset$. An analogous argument establishes that $Y_6 = \emptyset$.

Next, suppose that $Y_2 \neq \emptyset$, and fix $y_2 \in Y_2$. If y_0 is adjacent to y_2 , then y_0, x_1, x_2, y_2, y_0 is a 4-hole in G , contrary to the fact that G is C_4 -free. On the other hand, if y_0 is nonadjacent to y_2 , then $y_0, x_0, x_6, y_2, x_3, x_4, y_0$ is a 6-hole in G , contrary to the fact that G is C_6 -free. This proves that $Y_2 = \emptyset$. An analogous argument establishes that $Y_5 = \emptyset$.

Next, suppose that $Z_5 \neq \emptyset$, and fix $z_5 \in Z_5$. If y_0 is adjacent to z_5 , then y_0, x_4, x_5, z_5, y_0 is a 4-hole in G , contrary to the fact that G is C_4 -free. On the other hand, if y_0 is nonadjacent to z_5 , then $y_0, x_0, z_5, x_2, x_3, x_4, y_0$ is a 6-hole in G , contrary to the fact that G is C_6 -free. This proves that $Z_5 = \emptyset$. An analogous argument establishes that $Z_6 = \emptyset$.

It remains to show that at most one of Y_3, Y_4 is nonempty. But the proof of this is analogous to the above proof of the fact that, if $Y_0 \neq \emptyset$, then $Y_1 = \emptyset$. $\blacklozenge$

Claim 3.8.4. *For all $i \in \mathbb{Z}_7$, if $Z_i \neq \emptyset$, then Z_{i+2}, Z_{i+5} are empty.*

Proof of Claim 3.8.4. By symmetry, it suffices to prove the claim for $i = 0$. So, assume that $Z_0 \neq \emptyset$, and fix $z_0 \in Z_0$. Suppose that $Z_2 \neq \emptyset$, and fix $z_2 \in Z_2$. If z_0 is adjacent to z_2 , then z_0, z_2, x_6, x_0, z_0 is a 4-hole in G , contrary to the fact that G is C_4 -free. On the other hand, if z_0 is nonadjacent to z_2 , then z_0, x_2, z_2, x_4, z_0 is a 4-hole in G , contrary to the fact that G is C_4 -free. This proves that $Z_2 = \emptyset$. An analogous argument establishes that $Z_5 = \emptyset$. $\blacklozenge$

Claim 3.8.5. *For all $i \in \mathbb{Z}_7$, Y_i is complete to $Y_{i+3}, Y_{i+4}, Z_i, Z_{i+1}, Z_{i+3}, Z_{i+4}, W$ and anticomplete to Z_{i+2} .*

Proof of Claim 3.8.5. By symmetry, it suffices to prove the claim for $i = 0$. Note that:

- distinct, nonadjacent vertices x_0, x_4 are complete to $Y_0 \cup Y_3 \cup Z_0 \cup Z_3 \cup Z_4 \cup W$;
- distinct, nonadjacent vertices x_1, x_4 are complete to $Y_0 \cup Y_4 \cup Z_1$.

So, by Proposition 2.4, $Y_0 \cup Y_3 \cup Z_0 \cup Z_3 \cup Z_4 \cup W$ and $Y_0 \cup Y_4 \cup Z_1$ are both cliques, and it follows that Y_0 is complete to $Y_3, Y_4, Z_0, Z_1, Z_3, Z_4, W$. On the other hand, if some $y_0 \in Y_0$ and $z_2 \in Z_2$ are adjacent, then y_0, z_2, x_2, x_1, y_0 is a 4-hole in G , contrary to the fact that G is C_4 -free. It follows that Y_0 is anticomplete to Z_2 . $\blacklozenge$

Claim 3.8.6. *For all $i \in \mathbb{Z}_7$, Z_i is complete to $Z_{i+1}, Z_{i+3}, Z_{i+4}, Z_{i+6}, W$.*

Proof of Claim 3.8.6. By symmetry, it suffices to prove the claim for $i = 0$. Note that:

- distinct, nonadjacent vertices x_1, x_4 are complete to $Z_0 \cup Z_1 \cup Z_4 \cup W$;
- distinct, nonadjacent vertices x_0, x_3 are complete to $Z_0 \cup Z_3 \cup Z_6$.

So, by Proposition 2.4, $Z_0 \cup Z_1 \cup Z_4 \cup W$ and $Z_0 \cup Z_3 \cup Z_6$ are both cliques, and it follows that Z_0 is complete to Z_1, Z_3, Z_4, Z_6, W . $\blacklozenge$

Claim 3.8.7. *$(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of $G \setminus A$.*

Proof of Claim 3.8.7. Recall that sets $X_0, \dots, X_6, Y_0, \dots, Y_6, Z_0, \dots, Z_6, W$ are cliques that together partition $V(G) \setminus A$, and that cliques $X_0, \dots, X_6$ are all nonempty.²⁰ The result now follows immediately from the definition of a special partition, and from Claims 3.8.1-3.8.6. $\blacklozenge$

Claim 3.8.8. *A is anticomplete to X , and consequently, $N_G(A) \subseteq Y \cup Z \cup W$.*

Proof of Claim 3.8.8. Since $V(G) = X \cup Y \cup Z \cup W \cup A$, it suffices to show that A is anticomplete to X . Suppose otherwise. By symmetry, we may assume that some $a \in A$ and $x'_0 \in X_0$ are adjacent. But then a has exactly one neighbor in the 7-hole $x'_0, x_1, x_2, x_3, x_4, x_5, x_6, x'_0$, contrary to Lemma 3.7. $\blacklozenge$

Claim 3.8.9. *If $A \neq \emptyset$, and if B is the vertex set of a component of $G[A]$, then $N_G(B)$ is a clique-cutset of G .*

Proof of Claim 3.8.9. Assume that $A \neq \emptyset$, and let B be the vertex set of some component of $G[A]$. By Claim 3.8.8, B is anticomplete to X , and we have that $N_G(B) \subseteq Y \cup Z \cup W$. In particular, $N_G(B)$ is a cutset of G that separates B from X in G . It remains to show that $N_G(B)$ is a clique. Suppose otherwise. It then follows from the definition of a special partition (or alternatively, from Lemma 3.1) that there exists some $i \in \mathbb{Z}_7$ such that $N_G(B)$ intersects both Y_i and Z_{i+2} . By symmetry, we may assume that $N_G(B)$ intersects both Y_0 and Z_2 . Fix vertices $y_0 \in Y_0$ and $z_2 \in Z_2$ that both have a neighbor in B . Since $G[B]$ is connected, there exists an induced path $b_0, \dots, b_t$ ($t \geq 0$) in B such that b_0 is adjacent to y_0 , and b_t is adjacent to z_2 ; we may assume that $b_0, \dots, b_t$ is the shortest path with this property, so that y_0 is anticomplete to $\{b_1, \dots, b_t\}$, and z_2 is anticomplete to $\{b_0, \dots, b_{t-1}\}$. Clearly, $\{b_0, \dots, b_t\}$ is anticomplete to X (because B is anticomplete to X). But now if t is odd, then $y_0, b_0, \dots, b_t, z_2, x_2, x_1, y_0$ is an even hole in G ; and if t is even, then $y_0, b_0, \dots, b_t, z_2, x_4, y_0$ is an even hole in G . In either case, G contains an even hole, contrary to the fact that G is (P_7, C_4, C_6) -free, and therefore even-hole-free. $\blacklozenge$

Claim 3.8.10. *If G is $2P_3$ -free, then G is a 7-saucer with an associated 7-saucer partition $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$.*

Proof of Claim 3.8.10. Assume that G is $2P_3$ -free. First of all, by Claim 3.8.7, we know that $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of $G \setminus A$. So, if $A = \emptyset$, then G is a 7-saucer with an associated 7-saucer partition $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$, and we are done. We may therefore assume that $A \neq \emptyset$. By Claim 3.8.8, we know that $N_G(A) \subseteq Y \cup Z \cup W$.

First, we show that for all $i \in \mathbb{Z}_7$, either A is anticomplete to Y_i , or $Z_{i+2} = \emptyset$. Suppose otherwise. By symmetry, we may assume that A is not anticomplete to Y_0 and that $Z_2 \neq \emptyset$. Fix adjacent vertices $a \in A$ and $y_0 \in Y_0$, and fix any $z_2 \in Z_2$. Since $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ is a special partition of $G \setminus A$, we know that y_0 and z_2 are nonadjacent, and so by Claim 3.8.9, a is nonadjacent to z_2 .²¹ But now $G[a, y_0, x_0, x_2, z_2, x_5]$ is a $2P_3$, a contradiction.

Next, if a_0, a_1, a_2 is an induced two-edge path in $G[A]$, then $G[a_0, a_1, a_2, x_0, x_1, x_2]$ is a $2P_3$, a contradiction. This proves that $G[A]$ is P_3 -free. Now Lemma 2.5 guarantees that A can be partitioned into nonempty cliques, say $A_1, \dots, A_\ell$, pairwise anticomplete to each other.

²⁰ Cliques $Y_0, \dots, Y_6, Z_0, \dots, Z_6, W$ may possibly be empty.

²¹ Indeed, let B be the vertex set of a component of $G[A]$ such that $a \in B$. By Claim 3.8.9, $N_G(B)$ is a clique, and so since y_0, z_2 are nonadjacent, $N_G(B)$ may contain at most one of them. So, a can be adjacent to at most one of y_0, z_2 . Since a is in fact adjacent to y_0 , it follows that a is nonadjacent to z_2 .

Finally, fix any $i \in \{1, \dots, \ell\}$. By Claim 3.8.9, $N_G(A_i)$ is a clique. Now Proposition 2.6 guarantees that A_i can be ordered as $A_i = \{a_1^i, \dots, a_{r_i}^i\}$ so that $N_G[a_{r_i}^i] \subseteq \dots \subseteq N_G[a_1^i]$.²²

We have now shown that G is a 7-saucer, and that $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$ is an associated 7-saucer partition of G . ♦

We are now done: Claims 3.8.7 and 3.8.9 together guarantee that G admits either a special partition or a clique-cutset, whereas Claim 3.8.10 guarantees that if G is $2P_3$ -free, then G is a 7-saucer. □

Theorem 3.9. *For any graph G , the following are equivalent:*

- G is $(2P_3, C_4, C_6)$ -free and contains an induced C_7 ;
- G is a 7-saucer.

Proof. This follows immediately from Proposition 3.4 and Lemma 3.8. □

Lemma 3.10. *If a $4K_1$ -free graph G contains an induced C_7 , then G does not admit a clique-cutset.*

Proof. Fix a $4K_1$ -free graph G that contains a 7-hole (say, $x_0, x_1, x_2, \dots, x_6, x_0$), and suppose that G admits a clique-cutset C . Let (A, B) be a partition of $V(G) \setminus C$ into nonempty sets, anticomplete to each other. Since C_7 does not admit a clique-cutset, we see that $X := \{x_0, x_1, \dots, x_6\}$ intersects at most one of A and B ; by symmetry, we may assume that $X \cap B = \emptyset$, and consequently, $X \subseteq A \cup C$. But since C is a clique, we see that $G[X \setminus C]$ contains an induced P_5 , and consequently, an induced $3K_1$. The three vertices of this $3K_1$ together with any vertex of B induce a $4K_1$ in G , a contradiction. □

Theorem 3.11. *For any graph G , the following are equivalent:*

- (a) G is a (P_7, C_4, C_6) -free graph that contains an induced C_7 and does not admit a clique-cutset;
- (b) G is a $(2P_3, C_4, C_6)$ -free graph that contains an induced C_7 and contains no simplicial vertices;
- (c) G is a $(4K_1, C_4, C_6)$ -free graph that contains an induced C_7 ;
- (d) G admits a special partition;
- (e) G has exactly one nontrivial anticomponent, and this anticomponent is a thickening of a graph in $\mathcal{M}$;
- (f) G can be obtained from a thickening of a graph in $\mathcal{M}$ by possibly adding universal vertices to it.

Proof. By Proposition 3.6, (d), (e), and (f) are equivalent. By Lemma 3.8, (a) implies (d). Further, by Lemma 3.2, (d) implies (c). It remains to show that (c) implies (b), and that (b) implies (a).

We first assume (c) and prove (b). Since G is $(4K_1, C_4, C_6)$ -free, and since $4K_1$ is an induced subgraph of $2P_3$, we know that G is $(2P_3, C_4, C_6)$ -free. Moreover, by Lemma 3.10, G does not

²²Indeed, Proposition 2.6(b) applied to the cliques A_i and $N_G(A_i)$ guarantees that A_i can be ordered as $A_i = \{a_1^i, \dots, a_{r_i}^i\}$ so that $N_G(a_{r_i}^i) \cap N_G(A_i) \subseteq \dots \subseteq N_G(a_1^i) \cap N_G(A_i)$. Since A_i is a clique, this implies $N_G[a_{r_i}^i] \subseteq \dots \subseteq N_G[a_1^i]$.

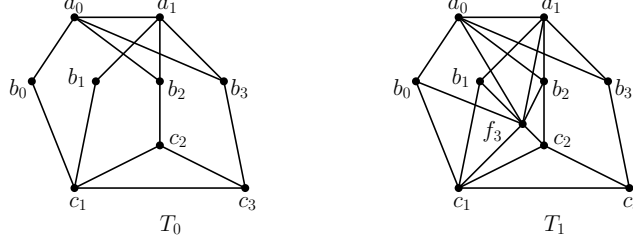


Figure 4.1: Graphs T_0 (left) and T_1 (right).

admit a clique-cutset; since the graph G is not complete (because it contains an induced C_7), it follows that G contains no simplicial vertices.²³ This proves (b).

To complete the proof, it remains to show that (b) implies (a). Since $2P_3$ is an induced subgraph of P_7 , it is clear that any $(2P_3, C_4, C_6)$ -free graph that contains an induced C_7 is, in particular, a (P_7, C_4, C_6) -free graph that contains an induced C_7 . This, together with Proposition 2.7, guarantees that (b) implies (a). $\square$

4 $(2P_3, C_4, C_6, C_7)$ -free graphs that contain an induced T_0

In this section, we give a full structural description of all $(2P_3, C_4, C_6, C_7)$ -free graphs that contain an induced T_0 (see Theorem 4.6). As an easy corollary, we obtain a full structural description of $(2P_3, C_4, C_6, C_7)$ -free graphs that contain an induced T_0 and contain no simplicial vertices (see Corollary 4.7).

We begin with some definitions. First of all, the *3-pentagon*, as well as graphs T_0 and T_1 , are represented in Figure 1.2. Note that the 3-pentagon is an induced subgraph of T_0 , and T_0 is an induced subgraph of T_1 . Next, recall that a graph is *chordal* if it contains no holes. A *simplicial elimination ordering* of a graph G is an ordering $v_1, \dots, v_t$ of the vertices of G such that for all $i \in \{1, \dots, t\}$, the vertex v_i is simplicial in the graph $G \setminus \{v_1, \dots, v_{i-1}\}$. It is well known that a graph is chordal if and only if it admits a simplicial elimination ordering [16].

Proposition 4.1. *T_0 and T_1 are both anticonnected, contain no simplicial vertices and no universal vertices, and are $(2P_3, C_4, C_6, C_7)$ -free. Consequently, any thickening of T_0 or T_1 is anticonnected, contains no simplicial and no universal vertices, and is $(2P_3, C_4, C_6, C_7)$ -free.*

Proof. In view of Proposition 2.10, it is enough to prove that T_0 and T_1 are both anticonnected, contain no simplicial vertices and no universal vertices, and are $(2P_3, C_4, C_6, C_7)$ -free. We will use the vertex labeling of T_0 and T_1 from Figure 4.1, so that $T_1 \setminus f_3 = T_0$. In both T_0 and T_1 , vertices b_0, b_1, b_2, b_3 together form a stable set of size four, and all the other vertices have at least one nonneighbor in this stable set; therefore, T_0 and T_1 are both anticonnected. The fact that T_0 and T_1 contain no simplicial vertices and no universal vertices can easily be seen by inspection. It remains to show that T_0 and T_1 are $(2P_3, C_4, C_6, C_7)$ -free. Since T_0 is an induced subgraph of T_1 , it in fact suffices to show that T_1 is $(2P_3, C_4, C_6, C_7)$ -free.

Claim 4.1.1. *All holes in T_1 are of length five. Consequently, T_1 is (C_4, C_6, C_7) -free.*

²³This is because the neighborhood of any simplicial vertex in a graph that is not complete is a clique-cutset of that graph.

Proof of Claim 4.1.1. Suppose that $x_0, x_1, \dots, x_{k-1}, x_0$ (where $k \geq 4$ and the indices are understood to be in $\mathbb{Z}_k$) is a hole in T_1 , and set $X := \{x_0, x_1, \dots, x_{k-1}\}$. We must show that $k = 5$.

Suppose first that $f_3 \in X$. By symmetry, we may assume that $x_0 = f_3$. Since $V(T_1) \setminus N_{T_1}[f_3] = \{b_3, c_3\}$, and since the nonneighbors of $x_0 = f_3$ in the hole $x_0, x_1, \dots, x_{k-1}, x_0$ are precisely the vertices $x_2, \dots, x_{k-2}$, we deduce that $x_2, \dots, x_{k-2} \in \{b_3, c_3\}$. Consequently, $4 \leq k \leq 5$. If $k = 5$, then we are done. We may therefore assume that $k = 4$ and $x_2 \in \{b_3, c_3\}$. Now, note that $x_0 = f_3$ and x_2 have two nonadjacent common neighbors (namely x_1 and x_3) in the hole x_0, x_1, x_2, x_3, x_0 . But the two common neighbors of f_3 and b_3 (namely a_0 and a_1) are adjacent, as are the two common neighbors of f_3 and c_3 (namely c_1 and c_2). Thus, $x_2 \neq b_3$ and $x_2 \neq c_3$, contrary to the fact that $x_2 \in \{b_3, c_3\}$.

From now on, we may assume that $f_3 \notin X$, which means that our hole $x_0, x_1, \dots, x_{k-1}, x_0$ is in fact a hole in $T_1 \setminus f_3 = T_0$. First of all, note that $b_0, b_1, b_2, b_3, c_1, c_2, c_3$ is a simplicial elimination ordering of $T_0 \setminus \{a_0, a_1\}$, and so by [16], $T_0 \setminus \{a_0, a_1\}$ is chordal, i.e. it contains no holes. So, at last one of a_0, a_1 belongs to X .

Suppose first that $a_0, a_1 \in X$. Since a_0 and a_1 are adjacent, we may assume by symmetry that $x_0 = a_0$ and $x_1 = a_1$. Then $x_{k-1} \in N_{T_0}(a_0) \setminus N_{T_0}(a_1)$ and $x_2 \in N_{T_0}(a_1) \setminus N_{T_0}(a_0)$, and we deduce that $x_{k-1} = b_0$ and $x_2 = b_1$. Then $x_{k-2} \in N_{T_0}(x_{k-1}) \setminus N_{T_0}(x_0) = N_{T_0}(b_0) \setminus N_{T_0}(a_0) = \{c_1\}$ and $x_3 \in N_{T_0}(x_2) \setminus N_{T_0}(x_1) = N_{T_0}(b_1) \setminus N_{T_0}(a_1) = \{c_1\}$. Thus, $x_{k-2} = x_3 = c_1$; we now deduce that $k - 2 = 3$, and therefore, $k = 5$.

From now on, we may assume that exactly one of a_0, a_1 belongs to X . By symmetry, we may assume that $a_1 \in X$ and $a_0 \notin X$. Consequently, $x_0, x_1, \dots, x_{k-1}, x_0$ is a hole in $T_0 \setminus a_0$. Since $d_{T_0 \setminus a_0}(b_0) = 1$, we see that $b_0 \notin X$. Therefore, $x_0, x_1, \dots, x_{k-1}, x_0$ is a hole in the 3-pentagon $T_0 \setminus \{a_0, b_0\}$, and it is clear that all holes in a 3-pentagon are of length five. Therefore, $k = 5$, and we are done. $\blacklozenge$

Claim 4.1.2. T_1 is $2P_3$ -free.

Proof of Claim 4.1.2. Since $V(T_1) \setminus N_{T_1}[f_3] = \{b_2, c_3\}$ is a clique, Proposition 2.8 guarantees that T_1 is $2P_3$ -free if and only if $T_1 \setminus f_3 = T_0$ is a $2P_3$ -free. Therefore, it suffices to show that T_0 is $2P_3$ -free. Suppose otherwise, and fix pairwise distinct vertices $x_0, x_1, x_2, y_0, y_1, y_2 \in V(T_0)$ such that x_0, x_1, x_2 and y_0, y_1, y_2 are both induced paths in T_0 , and such that $\{x_0, x_1, x_2\}$ is anticomplete to $\{y_0, y_1, y_2\}$.

Suppose first that $a_0, a_1 \notin \{x_0, x_1, x_2, y_0, y_1, y_2\}$. Then vertices $x_0, x_1, x_2, y_0, y_1, y_2$ induce a $2P_3$ in the graph $T_0 \setminus \{a_0, a_1\}$. But note that, in the graph $T_0 \setminus \{a_0, a_1\}$, vertices b_0, b_1, b_2, b_3 have degree 1, whereas vertices x_1, y_1 have degree at least two. So, $x_1, y_1 \notin \{b_0, b_1, b_2, b_3\}$, and therefore, $x_1, y_1 \in \{c_1, c_2, c_3\}$. But this is impossible since x_1, y_1 are nonadjacent, whereas $\{c_1, c_1, c_3\}$ is a clique.

Suppose next that $a_0, a_1 \in \{x_0, x_1, x_2, y_0, y_1, y_2\}$. Since a_0, a_1 are adjacent, we may assume by symmetry that $x_0 = a_0$ and $x_1 = a_1$. Then $x_2 \in N_{T_0}(x_1) \setminus N_{T_0}[x_0] = N_{T_0}(a_1) \setminus N_{T_0}[a_0] = \{b_1\}$. So, $x_2 = b_1$. But now $y_0, y_1, y_2 \in V(T_0) \setminus N_{T_0}[\{x_0, x_1, x_2\}] = V(T_0) \setminus N_{T_0}[\{a_0, a_1, b_1\}] = \{c_2, c_3\}$. Since y_0, y_1, y_2 are pairwise distinct, this is clearly impossible.

Thus, $\{x_0, x_1, x_2, y_0, y_1, y_2\}$ contains exactly one of a_0, a_1 . By symmetry, we may assume that $a_1 \in \{x_0, x_1, x_2, y_0, y_1, y_2\}$ and $a_0 \notin \{x_0, x_1, x_2, y_0, y_1, y_2\}$, so that vertices $x_0, x_1, x_2, y_0, y_1, y_2$ together induce a $2P_3$ in the graph $T_0 \setminus a_0$.

Suppose first that $b_0 \in \{x_0, x_1, x_2, y_0, y_1, y_2\}$. Since $d_{T_0 \setminus a_0}(b_0) = 1$, we may assume by symmetry that $b_0 = x_0$. But then $x_1 = c_1$ and $x_2 \in \{b_1, c_2, c_3\}$.²⁴ In any case, we have that

²⁴The fact that $x_1 = c_1$ follows from the fact that c_1 is the only neighbor of $x_1 = b_0$ in $T_0 \setminus a_0$. On the other hand, the fact that $x_2 \in \{b_1, c_2, c_3\}$ follows from the fact that $N_{T_0 \setminus a_0}(c_1) = \{b_0, b_1, c_2, c_3\}$ and the fact that

$a_1 \in \{y_0, y_1, y_2\}$, and consequently, a_1 is nonadjacent to x_2 . Therefore, $x_2 \neq b_1$, and it follows that $x_2 \in \{c_2, c_3\}$. By symmetry, we may assume that $x_2 = c_2$. But now y_0, y_2, y_3 is an induced two-edge-path in the graph $T_0 \setminus N_{T_0}[\{x_0, x_1, x_2\}] = T_0 \setminus N_{T_0}[\{b_0, c_1, c_2\}] = T_0[a_1, b_3]$, which is impossible, since this graph contains only two vertices.

Thus, $b_0 \notin \{x_0, x_1, x_2, y_0, y_1, y_2\}$. It follows that $T_0 \setminus \{a_0, b_0\}$ contains an induced $2P_3$. But note that $T_0 \setminus \{a_0, b_0\}$ is a 3-pentagon, and in particular, it contains exactly seven vertices. Therefore, by deleting one suitably chosen vertex of the 3-pentagon $T_0 \setminus \{a_0, b_0\}$, we should obtain a $2P_3$. But it is easy to see by inspection that no such vertex exists. This proves that T_0 is, in fact, $2P_3$ -free, and we are done. $\blacklozenge$

By Claims 4.1.1 and 4.1.2, we have that T_1 is $(2P_3, C_4, C_6, C_7)$ -free, and we are done. $\square$

Lemma 4.2. *Let G be a $(2P_3, C_4, C_6)$ -free graph, and let T an induced subgraph of G that is isomorphic to T_0 , with the vertices of T labeled as in Figure 4.1 (left). Then for all $x \in V(G)$, one of the following holds:*

- (a) *there exists some $v \in V(T)$ such that $N_G[x] \cap V(T) = N_T[v]$;*
- (b) *there exists some $i \in \{2, 3\}$ such that $N_G[x] \cap V(T) = V(T) \setminus \{b_i, c_i\}$;*
- (c) *$N_G[x] \cap V(T) = \{c_2, c_3\}$;*
- (d) *$N_G[x] \cap V(T) = \emptyset$;*
- (e) *$N_G[x] \cap V(T) = V(T)$.*

Proof. Fix an arbitrary $x \in V(G)$. If $x \in V(T)$, then clearly, $N_G[x] \cap V(T) = N_T[x]$, and in particular, (a) holds. We may therefore assume that $x \notin V(T)$. We now prove a sequence of claims.

Claim 4.2.1. *If x has a neighbor both in $\{b_0, b_1\}$ and in $\{b_2, b_3\}$, then x is complete to $\{a_0, a_1\}$ and has at least three neighbors in $\{b_0, b_1, b_2, b_3\}$.*

Proof of Claim 4.2.1. Assume that x has a neighbor both in $\{b_0, b_1\}$ and in $\{b_2, b_3\}$; by symmetry, we may assume that x is adjacent to b_1 and b_2 . Then x is adjacent to a_1 , for otherwise, x, b_1, a_1, b_2, x would be a 4-hole in G , a contradiction. It now remains to show that x has a neighbor in $\{b_0, b_3\}$ and is adjacent to a_0 .

Suppose that x is anticomplete to $\{b_0, b_3\}$. Then x is nonadjacent to c_3 , for otherwise, x, c_3, b_3, a_1, x would be a 4-hole in G , a contradiction. Then x has a neighbor in $\{c_1, c_2\}$, for otherwise, $G[x, a_1, b_3, b_0, c_1, c_2]$ would be a $2P_3$, a contradiction. If x is adjacent to c_1 , but nonadjacent to c_2 , then x, c_1, c_2, b_2, x is a 4-hole in G , a contradiction. Similarly, if x is adjacent to c_2 , but nonadjacent to c_1 , then x, c_2, c_1, b_1, x is a 4-hole in G , a contradiction. So, x is complete to $\{c_1, c_2\}$. Then x is nonadjacent to a_0 , for otherwise, x, a_0, b_0, c_1, x would be a 4-hole in G , a contradiction. But now $x, c_1, c_3, b_3, a_0, b_2, x$ is a 6-hole in G , a contradiction.

We have now shown that x has a neighbor in $\{b_0, b_3\}$, and it remains to show that x is adjacent to a_0 . Fix $i \in \{0, 3\}$ such that x is adjacent to b_i . Then x is adjacent to a_0 , for otherwise, x, b_i, a_0, b_2, x would be a 4-hole in G , a contradiction. $\blacklozenge$

Claim 4.2.2. *If x has a neighbor in both $\{b_0, b_1\}$ and $\{b_2, b_3\}$, then one of (a), (b), (e) holds.*

$b_0 = x_0$.

Proof of Claim 4.2.2. By symmetry, we may assume that x is adjacent to b_1 and b_2 . By Claim 4.2.1, x is complete to $\{a_0, a_1\}$ and is adjacent to at least one of b_0, b_3 .

Suppose first that x is adjacent to b_0 . Then x is adjacent to c_1 , for otherwise, x, b_0, c_1, b_1, x would be a 4-hole, a contradiction. Further, x is adjacent to c_2 , for otherwise, x, c_1, c_2, b_2, x would be a 4-hole, a contradiction. If x is complete to $\{b_3, c_3\}$, then (e) holds, and if x is anticomplete to $\{b_3, c_3\}$, then (b) holds (with $i = 3$). So, we may assume that x is mixed on $\{b_3, c_3\}$. But if x is adjacent to b_3 and nonadjacent to c_3 , then x, c_2, c_3, b_3, x is a 4-hole in G , a contradiction; and if x is adjacent to c_3 and nonadjacent to b_3 , then x, a_1, b_3, c_3, x is a 4-hole in G , a contradiction.

From now on, we assume that x is nonadjacent to b_0 , and is consequently adjacent to b_3 . Then x is nonadjacent to c_1 , for otherwise, x, c_1, b_0, a_0, x would be a 4-hole in G , a contradiction. But then x is in fact anticomplete to $\{c_1, c_2, c_3\}$, for otherwise, we fix an index $i \in \{2, 3\}$ such that x is adjacent to c_i , and we observe that x, c_i, c_1, b_1, x is a 4-hole in G . We now see that $N_G(x) \cap (T) = N_G[a_1]$, and in particular, (a) holds. $\blacklozenge$

Claim 4.2.3. *If x has at least two neighbors in $\{b_0, b_1, b_2, b_3\}$, then one of (a), (b), (e) holds.*

Proof of Claim 4.2.3. In view of Claim 4.2.2, we may assume that x is complete to one of $\{b_0, b_1\}$ and $\{b_2, b_3\}$ and is anticomplete to the other.

Suppose first that x is complete to $\{b_2, b_3\}$ and anticomplete to $\{b_0, b_1\}$. Then x is adjacent to c_1 , for otherwise, $G[b_2, x, b_3, b_0, c_1, b_1]$ would be a $2P_3$, a contradiction. But then if x is adjacent to a_0 , then x, a_0, b_0, c_1, x is a 4-hole in G , and if x is nonadjacent to a_0 , then x, b_2, a_0, b_3, x is a 4-hole in G , a contradiction in either case.

Thus, x is complete to $\{b_0, b_1\}$ and anticomplete to $\{b_2, b_3\}$. Then x is adjacent to c_1 , for otherwise, x, b_0, c_1, b_1, x would be a 4-hole in G , a contradiction. Further, x is adjacent to at least one of c_2, c_3 , for otherwise, $G[b_0, x, b_1, b_2, c_2, c_3]$ would be a $2P_3$, a contradiction. By symmetry, we may assume that x is adjacent to c_2 . Then x is anticomplete to $\{a_0, a_1\}$, for otherwise, we fix an index $i \in \{0, 1\}$ such that x is adjacent to a_i , and we observe that x, a_i, b_2, c_2, x is a 4-hole in G , a contradiction. Now x is adjacent to c_3 , for otherwise, $x, c_2, c_3, b_3, a_0, b_0, x$ would be a 6-hole in G , a contradiction. But now $N_G(x) \cap V(P) = N_T[c_1]$, and in particular, (a) holds. $\blacklozenge$

Claim 4.2.4. *If x has a unique neighbor in $\{b_0, b_1, b_2, b_3\}$, then (a) holds.*

Proof of Claim 4.2.4. Suppose first that x is adjacent to exactly one of b_0, b_1 and is anticomplete to $\{b_2, b_3\}$. By symmetry, we may assume that x is adjacent to b_0 and is nonadjacent to b_1 .

Suppose that x is nonadjacent to c_1 . Then x is in fact anticomplete to $\{c_1, c_2, c_3\}$, for otherwise, we fix an index $i \in \{2, 3\}$ such that x is adjacent to c_i , and we observe that x, c_i, c_1, b_0, x is a 4-hole in G , a contradiction. Then x is nonadjacent to a_1 , for otherwise, $x, a_1, b_2, c_2, c_1, b_0, x$ would be a 6-hole in G , a contradiction. But now $G[x, b_0, c_1, b_2, a_1, b_3]$ is a $2P_3$, a contradiction. This proves that x is in fact adjacent to c_1 .

We now see that x is nonadjacent to a_1 , for otherwise, x, a_1, b_1, c_1, x would be a 4-hole in G , a contradiction. Moreover, x is adjacent to a_0 , for otherwise, $G[x, c_1, b_1, b_2, a_0, b_3]$ would be a $2P_3$, a contradiction. Finally, x is anticomplete to $\{c_2, c_3\}$, for otherwise, we fix an index $i \in \{2, 3\}$ such that x is adjacent to c_i , and we observe that x, c_i, b_i, a_0, x is a 4-hole in G , a contradiction. But now $N_G(x) \cap V(T) = N_T[b_0]$, and in particular, (a) holds.

From now on, we may assume that x is anticomplete to $\{b_0, b_1\}$, and consequently has a unique neighbor in $\{b_2, b_3\}$. By symmetry, we may assume that x is adjacent to b_3 and is nonadjacent to b_2 .

Suppose first that x is adjacent to c_1 . Then x is anticomplete to $\{a_0, a_1\}$, for otherwise,

we fix an index $i \in \{0, 1\}$ such that x is adjacent to a_i , and we observe that x, a_i, b_i, c_1, x is a 4-hole in G , a contradiction. Further, x is adjacent to c_3 , for otherwise, x, b_3, c_3, c_1, x would be a 4-hole in G , a contradiction. Finally, x is adjacent to c_2 , for otherwise, $x, c_1, c_2, b_2, a_1, b_3, x$ would be a 6-hole in G , a contradiction. We now have that $N_G(x) \cap V(T) = N_T[c_3]$, and in particular, (a) holds.

From now on, we may assume that x is nonadjacent to c_1 . Suppose first that x is adjacent to c_2 . Then x is nonadjacent to a_1 , for otherwise, x, a_1, b_2, c_2, x would be a 4-hole in G , a contradiction. But now $x, b_3, a_1, b_1, c_1, c_2, x$ is a 6-hole in G , a contradiction. This proves that x is nonadjacent to c_2 . Then x is adjacent to a_0 , for otherwise, $G[x, b_3, a_0, b_1, c_1, c_2]$ would be a $2P_3$, a contradiction. Similarly, x is adjacent to a_1 , for otherwise, $G[x, b_3, a_1, b_0, c_1, c_2]$ would be a $2P_3$, a contradiction. Finally, x is adjacent to c_3 , for otherwise, $G[x, a_1, b_2, b_0, c_1, c_3]$ would be a $2P_3$, a contradiction. But now $N_G(x) \cap V(T) = N_T[b_3]$, and in particular, (a) holds. $\blacklozenge$

Claim 4.2.5. *If x is anticomplete to $\{b_0, b_1, b_2, b_3\}$, then (c) or (d) holds.*

Proof of Claim 4.2.5. Assume that x is anticomplete to $\{b_0, b_1, b_2, b_3\}$. We first show that x is anticomplete to $\{a_0, a_1\}$. Suppose otherwise. By symmetry, we may assume that x is adjacent to a_1 . Then x is anticomplete to $\{c_1, c_2, c_3\}$, for otherwise, we fix an index $i \in \{1, 2, 3\}$ such that x is adjacent to c_i , and we observe that x, a_1, b_i, c_i, x is a 4-hole in G , a contradiction. But now $G[x, a_1, b_3, b_0, c_1, c_2]$ is a $2P_3$, a contradiction. This proves that x is indeed anticomplete to $\{a_0, a_1\}$.

We now have that $N_G(x) \cap V(T) \subseteq \{c_1, c_2, c_3\}$. Then x is nonadjacent to c_1 , for otherwise, $G[x, c_1, b_0, b_2, a_1, b_3]$ would be a $2P_3$, a contradiction. It follows that $N_G(x) \cap V(T) \subseteq \{c_2, c_3\}$, and it remains to show that x is either complete or anticomplete to $\{c_2, c_3\}$. Suppose otherwise. By symmetry, we may assume that x is adjacent to c_2 and nonadjacent to c_3 . But then $G[x, c_2, c_3, a_0, a_1, b_1]$ is a $2P_3$, a contradiction. $\blacklozenge$

The result now follows immediately from Claims 4.2.3, 4.2.4, and 4.2.5. $\square$

A *tent* is a graph G whose vertex set can be partitioned into (possibly empty) sets $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, F_2, F_3, W, Y, Z$ such that all the following hold:

- $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3$ are nonempty cliques;
- F_2, F_3, Y are cliques, and at most one of them is nonempty;
- W is a (possibly empty) clique;
- A_0 is complete to A_1 ;
- A_0 is complete to B_0, B_2, B_3 and anticomplete to B_1, C_1, C_2, C_3 ;
- A_1 is complete to B_1, B_2, B_3 and anticomplete to B_0, C_1, C_2, C_3 ;
- B_0, B_1, B_2, B_3 are pairwise anticomplete to each other;
- C_1, C_2, C_3 are pairwise complete to each other;
- C_1 is complete to B_0, B_1 and anticomplete to B_2, B_3 ;
- C_2 is complete to B_2 and anticomplete to B_0, B_1, B_3 ;
- C_3 is complete to B_3 and anticomplete to B_0, B_1, B_2 ;
- F_2 is complete to $A_0, A_1, B_0, B_1, B_3, C_1, C_3$ and anticomplete to B_2, C_2 ;

- F_3 is complete to $A_0, A_1, B_0, B_1, B_2, C_1, C_2$ and anticomplete to B_3, C_3 ;
- W is complete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, F_2, F_3$;
- Y is complete to C_2, C_3 and anticomplete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1$;
- Z is anticomplete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, Y$;²⁵
- Y can be ordered as $Y = \{y_1, \dots, y_t\}$ so that $N_G[y_t] \subseteq \dots \subseteq N_G[y_1]$;²⁶
- if $Z \neq \emptyset$, then Z can be partitioned into nonempty cliques $Z_1, \dots, Z_\ell$, pairwise anticomplete to each other, and moreover, for each $i \in \{1, \dots, \ell\}$, Z_i can be ordered as $Z_i = \{z_1^i, \dots, z_{t_i}^i\}$ so that $N_G[z_{t_i}^i] \subseteq \dots \subseteq N_G[z_1^i]$.

Under these conditions, we also say that $(A_0, A_1; B_0, B_1, B_2, B_3; C_1, C_2, C_3; F_2, F_3; W; Y; Z)$ is a *tent partition* of the tent G .

Note that any thickening of T_0 or T_1 is a tent.²⁷

Proposition 4.3. *Let G be a tent, and let $(A_0, A_1; B_0, B_1, B_2, B_3; C_1, C_2, C_3; F_2, F_3; W; Y; Z)$ be a tent partition of G . Then all the following hold:*

- (a) *if $Y \neq \emptyset$, then some vertex of Y is simplicial in G ;*
- (b) *if $Z \neq \emptyset$, then some vertex of Z is simplicial in G ;*
- (c) *every vertex of W is a universal vertex of $G \setminus (Y \cup Z)$;*
- (d) *$G \setminus (W \cup Y \cup Z)$ is a thickening of T_0 or T_1 .*

Proof. We first prove (a). Suppose that $Y \neq \emptyset$, and fix an ordering $Y = \{y_1, \dots, y_t\}$ of Y such that $N_G[y_t] \subseteq \dots \subseteq N_G[y_1]$, as in the definition of a tent. Our goal is to show that $N_G[y_t]$ is a clique, which will immediately imply that y_t is a simplicial vertex. Since $Y \neq \emptyset$, the definition of a tent guarantees that $F_2 = F_3 = \emptyset$, and it further guarantees that Y is a clique, complete to $C_2 \cup C_3$ and anticomplete to $A_0 \cup A_1 \cup B_0 \cup B_1 \cup B_2 \cup B_3 \cup C_1$. Therefore, $Y \cup C_2 \cup C_3 \subseteq N_G[y_t] \subseteq \dots \subseteq N_G[y_1] \subseteq Y \cup C_2 \cup C_3 \cup W$. Set $W_t := N_G(y_t) \cap W$, so that $N_G[y_t] = Y \cup C_2 \cup C_3 \cup W_t$. By the definition of a tent, $Y \cup C_2 \cup C_3$ is a clique, as is W_t (because it is a subset of the clique W). So, it suffices to show that $Y \cup C_2 \cup C_3$ is complete to W_t . By the definition of a tent, $C_2 \cup C_3$ is complete to W (and therefore to W_t as well). On the other hand, since $Y = \{y_1, \dots, y_t\}$ and $W_t \subseteq N_G[y_t] \subseteq \dots \subseteq N_G[y_1]$, we see that Y is complete to W_t . This proves that $Y \cup C_2 \cup C_3$ is indeed complete to W_t , and (a) follows.

Next, we prove (b). Suppose that $Z \neq \emptyset$, and let $(Z_1, \dots, Z_\ell)$ be a partition of Z into cliques, pairwise anticomplete to each other, as per the definition of a tent. Using the definition of a tent, we fix an ordering $Z_1 = \{z_1, \dots, z_t\}$ of Z_1 such that $N_G[z_t] \subseteq \dots \subseteq N_G[z_1]$. Our goal is to show that $N_G[z_t]$ is a clique, which will immediately imply that z_t is simplicial. By the definition of a tent, we have that $N_G(Z_1) \subseteq F_2 \cup F_3 \cup W$. Since Z_1 is a clique, we conclude that $Z_1 \subseteq N_G[z_t] \subseteq \dots \subseteq N_G[z_1] \subseteq Z_1 \cup F_2 \cup F_3 \cup W$. Set $W_t := N_G(z_t) \cap (F_2 \cup F_3 \cup W)$, so that $N_G[z_t] = Z_1 \cup W_t$. By the definition of a tent, $F_2 \cup F_3 \cup W$ is a clique,²⁸ and consequently, W_t

²⁵Note that this implies that $N_G(Z) \subseteq F_2 \cup F_3 \cup W$.

²⁶It is possible that $Y = \emptyset$; in this case, we simply have that $t = 0$. On the other hand, our definition implies that if $Y \neq \emptyset$, then $F_2 = F_3 = \emptyset$. Moreover, our definition implies that $Y \cup C_2 \cup C_3 \subseteq N_G[y] \subseteq Y \cup C_2 \cup C_3 \cup W$ for all $y \in Y$, and that Y and $C_2 \cup C_3 \cup W$ are both cliques.

²⁷This follows immediately from the appropriate definitions, as long as we keep in mind the vertex labelings of T_0 and T_1 from Figure 4.1.

²⁸Indeed, by the definition of a tent, F_2, F_3, W are all cliques, at most one of F_2, F_3 is nonempty, and W is complete to $F_2 \cup F_3$. Thus, $F_2 \cup F_3 \cup W$ is a clique.

is a clique. Finally, since $Z_1 = \{z_1, \dots, z_t\}$ and $W_t \subseteq N_G[z_t] \subseteq \dots \subseteq N_G[z_1]$, we see that Z_1 is complete to W_t . This proves that $N_G[z_t] = Z_1 \cup W_t$ is a clique, and we deduce that z_t is a simplicial vertex. This completes the proof of (b).

Further, (c) readily follows from the fact that, by the definition of a tent, W is a clique, complete to $A_0 \cup A_1 \cup B_0 \cup B_1 \cup B_2 \cup B_3 \cup C_1 \cup C_2 \cup C_3 \cup F_2 \cup F_3 = V(Q) \setminus (W \cup Y \cup Z)$.

It remains to prove (d). Set $H := G \setminus (W \cup Y \cup Z)$, so that $V(H) = A_0 \cup A_1 \cup B_0 \cup B_1 \cup B_2 \cup B_3 \cup C_1 \cup C_2 \cup C_3 \cup F_2 \cup F_3$. If $F_2 \cup F_3 = \emptyset$, then it is clear that H is a thickening of T_0 (this follows from the vertex labeling of T_0 from Figure 4.1 and from the definition of a tent). So, assume that $F_2 \cup F_3 \neq \emptyset$. By the definition of a tent, we see that exactly one of F_2, F_3 is nonempty. By symmetry, we may assume that $F_3 \neq \emptyset$ and $F_2 = \emptyset$ (otherwise, we simply swap the roles of B_2, C_2, F_2 and B_3, C_3, F_3 , respectively). But then H is a thickening of T_1 (this follows from the vertex labeling of T_1 from Figure 4.1 and from the definition of a tent). This proves (d). $\square$

Proposition 4.4. *Every tent is $(2P_3, C_4, C_6, C_7)$ -free and contains an induced T_0 .*

Proof. Let G be a tent, and let $(A_0, A_1; B_0, B_1, B_2, B_3; C_1, C_2, C_3; F_2, F_3; W; Y; Z)$ be a tent partition of G . By Proposition 4.3(d), $G \setminus (W \cup Y \cup Z)$ is a thickening of T_0 or T_1 , and in particular, it contains an induced T_0 or T_1 . Since T_0 is an induced subgraph of T_1 , we now deduce that G contains an induced T_0 . It remains to show that G is $(2P_3, C_4, C_6, C_7)$ -free.

Claim 4.4.1. *No vertex of $Y \cup Z$ belongs to any hole of G .*

Proof of Claim 4.4.1. Suppose otherwise, and let X be the vertex set of a hole in G such that $X \cap (Y \cup Z) \neq \emptyset$. Set $Y_X := Y \cap X$ and $Z_X := Z \cap X$; then $\emptyset \neq Y_X \cup Z_X \subseteq X$. Next, we set $G_X := G \setminus ((Y \setminus Y_X) \cup (Z \setminus Z_X))$, so that G_X is a tent with an associated tent partition $(A_0, A_1; B_0, B_1, B_2, B_3; C_1, C_2, C_3; F_2, F_3; W; Y_X; Z_X)$, and X induces a hole in G_X . Since $Y_X \cup Z_X \neq \emptyset$, Proposition 4.3(a-b) guarantees that some vertex of $Y_X \cup Z_X$ is simplicial in G_X . Since $Y_X \cup Z_X \subseteq X$, we deduce that some vertex of X is simplicial in G , which is impossible since no vertex of a hole is simplicial. $\blacklozenge$

Claim 4.4.2. *No vertex of W belongs to any hole of G .*

Proof of Claim 4.4.2. By Claim 4.4.1, no vertex of $Y \cup Z$ belongs to any hole of G . Therefore, it suffices to show that no hole of $G \setminus (Y \cup Z)$ belongs to W . But by Proposition 4.3(c), every vertex of W is a universal vertex of $G \setminus (Y \cup Z)$. Since holes contain no universal vertices, we deduce that no hole of $G \setminus (Y \cup Z)$ contains a vertex of W . $\blacklozenge$

Claim 4.4.3. *G is (C_4, C_6, C_7) -free.*

Proof of Claim 4.4.3. In view of Claims 4.4.1 and 4.4.2, it suffices to show that the graph $H := G \setminus (W \cup Y \cup Z)$ is (C_4, C_6, C_7) -free. By Proposition 4.3(d), H is a thickening of T_0 or T_1 , and so by Proposition 4.1, H is (C_4, C_6, C_7) -free. $\blacklozenge$

In view of Claim 4.4.3, it now remains to show that G is $2P_3$ -free. Suppose otherwise, and fix pairwise distinct vertices $x_0, x_1, x_2, y_0, y_1, y_2 \in V(G)$ such that x_0, x_1, x_2 and y_0, y_1, y_2 are induced two-edge paths in G , and such that $\{x_0, x_1, x_2\}$ and $\{y_0, y_1, y_2\}$ are pairwise anticomplete to each other in G .

Claim 4.4.4. $\{x_0, x_1, x_2, y_0, y_1, y_2\} \cap W = \emptyset$.

Proof of Claim 4.4.4. Suppose otherwise. By symmetry, we may assume that there exists some $i \in \{0, 1, 2\}$ such that $x_i \in W$. Proposition 4.3(c) then guarantees that x_i is a universal vertex

of $G \setminus (Y \cup Z)$. Since x_i is anticomplete to $\{y_0, y_1, y_2\}$, it follows that $y_0, y_1, y_2 \in Y \cup Z$. So, y_0, y_1, y_2 is an induced two-edge path in $G[Y \cup Z]$. But this is impossible since, by the definition of a tent, $Y \cup Z$ can be partitioned into cliques, pairwise anticomplete to each other.²⁹ $\blacklozenge$

Claim 4.4.5. $\{x_0, x_1, x_2, y_0, y_1, y_2\} \cap (F_2 \cup F_3) = \emptyset$.

Proof of Claim 4.4.5. Suppose otherwise. By symmetry, we may assume that there exists some $i \in \{0, 1, 2\}$ such that $x_i \in F_3$. In particular, $F_3 \neq \emptyset$, and so by the definition of a tent, $F_2 = Y = \emptyset$. Since x_i is anticomplete to $\{y_0, y_1, y_2\}$, we deduce that $y_0, y_1, y_2 \in V(G) \setminus N_G[x_i]$. Now, since $x_i \in F_3$, it follows from the definition of a tent that $A_0 \cup A_1 \cup B_0 \cup B_1 \cup B_2 \cup C_1 \cup C_2 \cup F_3 \cup W \subseteq N_G[x_i]$. Since $F_2 = Y = \emptyset$, we deduce that $V(G) \setminus N_G[x_i] \subseteq B_3 \cup C_3 \cup Z$. So, y_0, y_1, y_2 is an induced two-edge path in $G[B_3 \cup C_3 \cup Z]$. But this is impossible, since by the definition of a tent, $B_3 \cup C_3 \cup Z$ can be partitioned into cliques, pairwise anticomplete to each other.³⁰ $\blacklozenge$

Claim 4.4.6. $\{x_0, x_1, x_2, y_0, y_1, y_2\} \cap Z = \emptyset$.

Proof of Claim 4.4.6. Suppose otherwise. By symmetry, we may assume that $\{x_0, x_1, x_2\} \cap Z \neq \emptyset$. (In particular, $Z \neq \emptyset$.) By the definition of a tent, Z can be partitioned into cliques, pairwise anticomplete to each other; so, $G[Z]$ is $2P_3$ -free. Since $\{x_0, x_1, x_2\} \cap Z \neq \emptyset$, and since x_0, x_1, x_2 is an induced two-edge path in G , it follows that $\{x_0, x_1, x_2\} \cap N_G(Z) \neq \emptyset$. But since $N_G(Z) \subseteq F_2 \cup F_3 \cup W$ (by the definition of a tent), this contradicts Claims 4.4.4 and 4.4.5. $\blacklozenge$

Claim 4.4.7. $\{x_0, x_1, x_2, y_0, y_1, y_2\} \cap Y = \emptyset$.

Proof of Claim 4.4.7. By Claims 4.4.4, 4.4.5, and 4.4.6, we have that $\{x_0, x_1, x_2, y_0, y_1, y_2\} \cap (W \cup F_2 \cup F_3 \cup Z) = \emptyset$. So, we may assume that $F_2 = F_3 = W = Z = \emptyset$, for otherwise, instead of G , we consider the tent $G \setminus (F_2 \cup F_3 \cup W \cup Z)$, with an associated tent partition $(A_0, A_1; B_0, B_1, B_2, B_3; C_1, C_2, C_3; \emptyset; \emptyset; \emptyset; Y; \emptyset)$.

Now, suppose that $\{x_0, x_1, x_2, y_0, y_1, y_2\} \cap Y \neq \emptyset$. By symmetry, we may assume that $\{y_0, y_1, y_2\} \cap Y \neq \emptyset$. Consequently, $\{y_0, y_1, y_2\} \cap (Y \cup C_2 \cup C_3) \neq \emptyset$. Since $Y \cup C_2 \cup C_3$ is a clique (by the definition of a tent), and since y_0, y_1, y_2 is an induced two-edge path in G , it follows that $\{y_0, y_1, y_2\} \cap N_G(Y \cup C_2 \cup C_3) \neq \emptyset$. Since $F_2 = F_3 = W = \emptyset$, the definition of a tent implies that $N_G(Y \cup C_2 \cup C_3) = C_1 \cup B_2 \cup B_3$. We now have that $\{y_0, y_1, y_2\}$ intersects both Y and $C_1 \cup B_2 \cup B_3$; since y_1 is complete to $\{y_0, y_2\}$, whereas Y and $C_1 \cup B_2 \cup B_3$ are anticomplete to each other, we may now assume by symmetry that $y_0 \in Y$ and $y_2 \in C_1 \cup B_2 \cup B_3$, and consequently, $y_1 \in C_2 \cup C_3$. Since B_2 is anticomplete to C_3 , and since B_3 is anticomplete to C_2 , we may now assume by symmetry that $y_1 \in C_2$ and $y_2 \in B_2 \cup C_1$, so that one of the following holds:

- (1) $y_0 \in Y, y_1 \in C_2, y_2 \in B_2$;
- (2) $y_0 \in Y, y_1 \in C_2, y_2 \in C_1$.

In either case, we know that $x_0, x_1, x_2 \in V(G) \setminus N_G(\{y_0, y_1, y_2\})$. We now remind the reader that $F_2 = F_3 = W = Z = \emptyset$, a fact that we implicitly use below.

Suppose first that (1) holds. It then follows from the definition of a tent that $V(G) \setminus N_G(\{y_0, y_1, y_2\}) = B_0 \cup B_1 \cup B_3$. So, x_0, x_1, x_2 is an induced two-edge path in $G[B_0 \cup B_1 \cup B_2]$.

²⁹Indeed, by the definition of a tent, Y is a clique, anticomplete to Z . Moreover, Z is either empty or can be partitioned into nonempty cliques, pairwise anticomplete to each other.

³⁰Indeed, by the definition of a tent, B_3 and C_3 are cliques, complete to each other and anticomplete to Z . (So, $B_2 \cup C_3$ is a clique, anticomplete to Z .) Meanwhile, Z is either empty or can be partitioned into nonempty cliques, pairwise anticomplete to each other.

But this is impossible since, by the definition of a tent, B_0, B_1, B_2 are disjoint cliques, pairwise anticomplete to each other.

Suppose now that (2) holds. It then follows from the definition of a tent that $V(G) \setminus N_G(\{y_0, y_1, y_2\}) = A_0 \cup A_1 \cup B_3$. So, x_0, x_1, x_2 is an induced two-edge path in $G[A_0 \cup A_1 \cup B_3]$. But this is impossible since, by the definition of a tent, $A_0 \cup A_1 \cup B_3$ is a clique. $\blacklozenge$

By Claims 4.4.4, 4.4.6, and 4.4.7, we have that $\{x_0, x_1, x_2, y_0, y_1, y_2\} \cap (W \cup Y \cup Z) = \emptyset$. So, $G \setminus (W \cup Y \cup Z)$ is not $2P_3$ -free. But this contradicts Propositions 4.1 and 4.3(d): indeed, by Proposition 4.3(d), $G \setminus (W \cup Y \cup Z)$ is a thickening of T_0 or T_1 , and so by Proposition 4.1, $G \setminus (W \cup Y \cup Z)$ is $2P_3$ -free. $\square$

Lemma 4.5. *Let G be a $(2P_3, C_4, C_6, C_7)$ -free graph that contains an induced T_0 . Then G is a tent.*

Proof. Let T be an induced subgraph of G that is isomorphic to T_0 , with the vertices of T labeled as in Figure 4.1 (left). We now define sets $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, F_2, F_3, W, Y, Z$ as follows:

- for each $i \in \{0, 1\}$, set $A_i := \{x \in V(G) \mid N_G[x] \cap V(T) = N_T[a_i]\}$;
- for each $i \in \{0, 1, 2, 3\}$, set $B_i := \{x \in V(G) \mid N_G[x] \cap V(T) = N_T[b_i]\}$;
- for each $i \in \{1, 2, 3\}$, set $C_i := \{x \in V(G) \mid N_G[x] \cap V(T) = N_T[c_i]\}$;
- for each $i \in \{2, 3\}$, set $F_i := \{x \in V(G) \mid N_G[x] \cap V(T) = V(T) \setminus \{b_i, c_i\}\}$;
- set $W := \{x \in V(G) \mid N_G[x] \cap V(T) = V(T)\}$;
- set $Y := \{x \in V(G) \mid N_G[x] \cap V(T) = \{c_2, c_3\}\}$;
- set $Z := \{x \in V(G) \mid N_G[x] \cap V(T) = \emptyset\}$.

We now prove a sequence of claims, which together establish that G is a tent with an associated tent partition $(A_0, A_1; B_0, B_1, B_2, B_3; C_1, C_2, C_3; F_2, F_3; W; Y; Z)$.

Claim 4.5.1. *Sets $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, F_2, F_3, W, Y, Z$ form a partition of $V(G)$. Moreover, all the following hold:*

- for each $i \in \{0, 1\}$, $a_i \in A_i$;
- for each $i \in \{0, 1, 2, 3\}$, $b_i \in B_i$;
- for each $i \in \{1, 2, 3\}$, $c_i \in C_i$.

In particular, $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3$ are all nonempty.

Proof of Claim 4.5.1. The fact that sets $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, F_2, F_3, W, Y, Z$ form a partition of $V(G)$ follows immediately from Lemma 4.2. The rest is immediate from the definitions of the relevant sets. $\blacklozenge$

Claim 4.5.2. *Sets $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, F_2, F_3, W, Y$ are all cliques. Moreover, all the following hold:*

- at most one of F_2, F_3, Y is nonempty;
- A_0 is complete to A_1 ;
- W is complete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, F_2, F_3$.

Proof of Claim 4.5.2. First, with Proposition 2.4 in mind, we observe the following:

- distinct, nonadjacent vertices b_2, b_3 are both complete to $A_0 \cup A_1 \cup W$;
- distinct, nonadjacent vertices a_0, c_1 are both complete to $B_0 \cup W$;
- for all $i \in \{1, 2, 3\}$, distinct, nonadjacent a_1, c_i are both complete to $B_i \cup W$;
- distinct, nonadjacent vertices b_0, b_1 are both complete to $C_1 \cup W$;
- for all $i \in \{2, 3\}$, distinct, nonadjacent vertices b_i, c_1 are both complete to $C_i \cup W$;
- distinct, nonadjacent vertices a_1, c_1 are complete to $F_2 \cup F_3 \cup W$.

So, by Proposition 2.4, all the following are cliques: $A_0 \cup A_1 \cup W$, $B_0 \cup W$, $B_1 \cup W$, $B_2 \cup W$, $B_3 \cup W$, $C_1 \cup W$, $C_2 \cup W$, $C_3 \cup W$, $F_2 \cup F_3 \cup W$. Thus, sets $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, F_2, F_3, W$ are all cliques, A_0 is complete to A_1 , and W is complete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, F_2, F_3$. Furthermore, Y is clique, for if some distinct $y, y' \in Y$ were nonadjacent, then $G[y, c_3, y', b_1, a_1, b_2]$ would be a $2P_3$, a contradiction. It remains to show that at most one of F_2, F_3, Y is nonempty.

Suppose first that F_2 and F_3 are both nonempty. Fix $f_2 \in F_2$ and $f_3 \in F_3$. We saw above that $F_2 \cup F_3 \cup W$ is a clique, and consequently, f_2, f_3 are adjacent. But now f_2, c_3, c_2, f_3, f_2 is a 4-hole in G , a contradiction. This proves that at most one of F_2, F_3 is nonempty.

Next, suppose that F_2 and Y are both nonempty, and fix some $f_2 \in F_2$ and $y \in Y$. If y, f_2 are adjacent, then y, c_2, c_1, f_2, y is a 4-hole, a contradiction. On the other hand, if y, f_2 are nonadjacent, then $G[b_0, f_2, b_1, y, c_2, b_2]$ is a $2P_3$, again a contradiction. This proves that at most one of F_2, Y is empty. Similarly, at most one of F_3, Y is empty, and we are done. $\blacklozenge$

Claim 4.5.3. A_0 is complete to B_0, B_2, B_3 and anticomplete to B_1, C_1, C_2, C_3 .

Proof of Claim 4.5.3. If some $a'_0 \in A_0$ and $b'_0 \in B_0$ are nonadjacent, then $G[b_2, a'_0, b_3, b'_0, c_1, b_1]$ is a $2P_3$, a contradiction. Therefore, A_0 is complete to B_0 .

If some $a'_0 \in A_0$ and $b'_2 \in B_2$ are nonadjacent, then $a'_0, a_1, b'_2, c_2, c_1, b_0, a'_0$ is a 6-hole in G , a contradiction. Therefore, A_0 is complete to B_2 , and analogously, A_0 is complete to B_3 .

If some $a'_0 \in A_0$ and $b'_1 \in B_1$ are adjacent, then $a'_0, b'_1, c_1, b_0, a'_0$ is a 4-hole in G , a contradiction. Therefore, A_0 is anticomplete to B_1 .

If some $a'_0 \in A_0$ and $c'_1 \in C_1$ are adjacent, then $a'_0, c'_1, c_2, b_2, a'_0$ is a 4-hole in G , a contradiction. Therefore, A_0 is anticomplete to C_1 .

If some $a'_0 \in A_0$ and $c'_2 \in C_2$ are adjacent, then $a'_0, b_0, c_1, c'_2, a'_0$ is a 4-hole in G , a contradiction. Therefore, A_0 is anticomplete to C_2 , and analogously, A_0 is anticomplete to C_3 . $\blacklozenge$

Claim 4.5.4. A_1 is complete to B_1, B_2, B_3 and anticomplete to B_0, C_1, C_2, C_3 .

Proof of Claim 4.5.4. The proof of this claim is completely analogous to the proof of Claim 4.5.3 (we simply swap the roles of A_0 and A_1 , and we swap the roles of B_0 and B_1). $\blacklozenge$

Claim 4.5.5. B_0, B_1, B_2, B_3 are pairwise anticomplete to each other.

Proof of Claim 4.5.5. If some $b'_0 \in B_0$ and $b'_1 \in B_1$ are adjacent, then $b'_0, b'_1, a_1, a_0, b'_0$ is a 4-hole in G , a contradiction. Therefore, B_0 is anticomplete to B_1 .

If for some $i \in \{2, 3\}$, some $b'_0 \in B_0$ is adjacent to some $b'_i \in B_i$, then $b'_0, b'_i, c_i, c_1, b'_0$ is a 4-hole in G , a contradiction. Thus, B_0 is anticomplete to B_2, B_3 .

If for some distinct $i, j \in \{1, 2, 3\}$, some $b'_i \in B_i$ and $b'_j \in B_j$ are adjacent, then $b'_i, b'_j, c_j, c_i, b'_i$ is a 4-hole in G , a contradiction. Therefore, B_1, B_2, B_3 are pairwise anticomplete to each other. $\blacklozenge$

Claim 4.5.6. C_1, C_2, C_3 are pairwise complete to each other.

Proof of Claim 4.5.6. Suppose otherwise, and fix distinct $i, j \in \{1, 2, 3\}$ such that some $c'_i \in C_i$ and $c'_j \in C_j$ are nonadjacent to each other. Let k be the unique index in $\{1, 2, 3\} \setminus \{i, j\}$. Then $c'_i, c_k, c'_j, b_j, a_1, b_i, c'_i$ is a 6-hole in G , a contradiction. $\blacklozenge$

Claim 4.5.7. C_1 is complete to B_0, B_1 and anticomplete to B_2, B_3 .

Proof of Claim 4.5.7. If some $c'_1 \in C_1$ and $b'_0 \in B_0$ are nonadjacent, then $G[b'_0, a_0, b_2, b_1, c'_1, c_3]$ is a $2P_3$, a contradiction. Therefore, C_1 is complete to B_0 , and analogously, C_1 is complete to B_1 .

If some $c'_1 \in C_1$ and $b'_2 \in B_2$ are adjacent, then $c'_1, b'_2, a_1, b_1, c'_1$ is a 4-hole in G , a contradiction. Therefore, C_1 is anticomplete to B_2 , and analogously, C_1 is anticomplete to B_3 . $\blacklozenge$

Claim 4.5.8. C_2 is complete to B_2 and anticomplete to B_0, B_1, B_3 .

Proof of Claim 4.5.8. If some $c'_2 \in C_2$ and $b'_2 \in B_2$ are nonadjacent, then $G[b'_2, a_1, b_3, b_0, c_1, c'_2]$ is a $2P_3$, a contradiction. Therefore, C_2 is complete to B_2 .

If some $c'_2 \in C_2$ is adjacent to some $b'_0 \in B_0$, then $c'_2, b'_0, a_0, b_2, c'_2$ is a 4-hole in G , a contradiction. Therefore, C_2 is anticomplete to B_0 .

If for some $i \in \{1, 3\}$, some $c'_2 \in C_2$ is adjacent to some $b'_i \in B_i$, then $c'_2, b'_i, a_1, b_2, c'_2$ is a 4-hole in G , a contradiction. Therefore, C_2 is anticomplete to B_1, B_3 . $\blacklozenge$

Claim 4.5.9. C_3 is complete to B_3 and anticomplete to B_0, B_1, B_2 .

Proof of Claim 4.5.9. The proof of this claim is analogous to the proof of Claim 4.5.8. $\blacklozenge$

Claim 4.5.10. F_2 is complete to $A_0, A_1, B_0, B_1, B_3, C_1, C_3$ and anticomplete to B_2, C_2 .

Proof of Claim 4.5.10. First, with Proposition 2.4 in mind, we observe the following:

- for all $i \in \{0, 1\}$, distinct, nonadjacent vertices b_i, b_3 are complete to $F_2 \cup A_i$;
- distinct, nonadjacent vertices a_0, c_1 are complete to $F_2 \cup B_0$;
- for all $i \in \{1, 3\}$, distinct, nonadjacent vertices a_1, c_i are complete to $F_2 \cup B_i$;
- distinct, nonadjacent vertices b_1, c_3 are complete to $F_2 \cup C_1$;
- distinct, nonadjacent vertices b_3, c_1 are complete to $F_2 \cup C_3$.

Proposition 2.4 now guarantees that $F_2 \cup A_0, F_2 \cup A_1, F_2 \cup B_0, F_2 \cup B_1, F_2 \cup B_3, F_2 \cup C_1, F_2 \cup C_3$ are all cliques. Consequently, F_2 is complete to $A_0, A_1, B_0, B_1, B_3, C_1, C_3$.

Next, if some $f_2 \in F_2$ and $b'_2 \in B_2$ are adjacent, then f_2, b'_2, c_2, c_3, f_2 is a 4-hole in G , a contradiction. Thus, F_2 is anticomplete to B_2 .

Finally, if some $f_2 \in F_2$ and $c'_2 \in C_2$ are adjacent, then f_2, c'_2, b_2, a_1, f_2 is a 4-hole in G , a contradiction. Thus, F_2 is anticomplete to C_2 . $\blacklozenge$

Claim 4.5.11. F_3 is complete to $A_0, A_1, B_0, B_1, B_2, C_1, C_2$ and anticomplete to B_3, C_3 .

Proof of Claim 4.5.11. The proof of this claim is analogous to the proof of Claim 4.5.10. $\blacklozenge$

Claim 4.5.12. *Y is complete to C_2, C_3 and anticomplete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1$.*

Proof of Claim 4.5.12. If some $y \in Y$ is nonadjacent to some $c'_2 \in C_2$, then we observe that $T' := G[(V(T) \setminus \{c_2\}) \cup \{c'_2\}]$ is a T_0 , and y has a unique neighbor in $V(T')$, contrary to Lemma 4.2. Therefore, Y is complete to C_2 , and analogously, Y is complete to C_3 .

It remains to show that Y is anticomplete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1$. Suppose otherwise, and fix $X \in \{A_0, A_1, B_0, B_1, B_2, B_3, C_1\}$ such that some $y \in Y$ and $x' \in X$ are adjacent. Let x be the unique vertex in $X \cap V(T)$. Then $T' := G[(V(T) \setminus \{x\}) \cup \{x'\}]$ is a T_0 , and y has exactly three neighbors (namely, x', c_2, c_3) in it. It then follows from Lemma 4.2 that there exists some $v \in V(T')$ such that $N_G(y) \cap V(T') = N_{T'}[v]$. But now $N_{T'}[v] = \{x', c_2, c_3\}$, and consequently, $v \in \{x', c_2, c_3\}$. Since $d_{T'}(c_2) = d_{T'}(c_3) = 3$, we deduce that $v = x'$, and consequently, $N_{T'}(x') = \{c_2, c_3\}$. But then $N_T(x) = \{c_2, c_3\}$, which is impossible since no vertex of T has neighborhood (in T) precisely $\{c_2, c_3\}$. ♦

Claim 4.5.13. *Z is anticomplete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, Y$.*

Proof of Claim 4.5.13. We first show that Z is anticomplete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3$. Suppose otherwise, and fix $X \in \{A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3\}$ such that some $z \in Z$ and $x' \in X$ are adjacent. Let x be the unique vertex in $V(T) \cap X$. Then $T' := G[(V(T) \setminus \{x\}) \cup \{x'\}]$ is a T_0 , and z has a unique neighbor in $V(T')$, contrary to Lemma 4.2. This proves that Z is indeed anticomplete to $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3$.

It remains to show that Z is anticomplete to Y . Suppose otherwise, and fix adjacent $z \in Z$ and $y \in Y$. But then $G[z, y, c_3, b_0, a_0, b_2]$ is a $2P_3$, a contradiction. ♦

Claim 4.5.14. *Y can be ordered as $Y = \{y_1, \dots, y_t\}$ so that $N_G[y_t] \subseteq \dots \subseteq N_G[y_1]$.*

Proof of Claim 4.5.14. We may assume that $Y \neq \emptyset$, for otherwise, the result is immediate (with $t = 0$). By Claim 4.5.2, Y is a clique, and $F_2 = F_3 = \emptyset$. So, by Claim 4.5.1, we have that sets $A_0, A_1, B_0, B_1, B_2, B_3, C_1, C_2, C_3, W, Y, Z$ form a partition of $V(G)$; this, together with Claims 4.5.12 and 4.5.13, and together with the fact that Y is a clique, implies that for all $y \in Y$, we have that $C_2 \cup C_3 \cup Y \subseteq N_G[y] \subseteq C_2 \cup C_3 \cup W \cup Y$. Therefore, it suffices to show that Y can be ordered as $Y = \{y_1, \dots, y_t\}$ so that $N_G(y_t) \cap W \subseteq \dots \subseteq N_G(y_1) \cap W$. But since Y and W are both cliques (by Claim 4.5.2), this follows immediately from Proposition 2.6. ♦

Claim 4.5.15. *If $Z \neq \emptyset$, then Z can be partitioned into nonempty cliques $Z_1, \dots, Z_\ell$, pairwise anticomplete to each other, and moreover, for each $i \in \{1, \dots, \ell\}$, Z_i can be ordered as $Z_i = \{z_1^i, \dots, z_{t_i}^i\}$ so that $N_G[z_{t_i}^i] \subseteq \dots \subseteq N_G[z_1^i]$.*

Proof of Claim 4.5.15. Assume that $Z \neq \emptyset$. First, since vertices b_1, a_1, b_2 induce a P_3 in G , and since they are all anticomplete to Z , we see that $G[Z]$ is P_3 -free (otherwise, G would contain an induced $2P_3$, a contradiction). Therefore, by Proposition 2.5, every component of $G[Z]$ is a complete graph, or in other words, Z can be partitioned into nonempty cliques $Z_1, \dots, Z_\ell$, pairwise anticomplete to each other.

Now, fix $i \in \{1, \dots, \ell\}$. By Claims 4.5.1 and 4.5.13, we see that every vertex $z \in Z_i$ satisfies $Z_i \subseteq N_G[z] \subseteq F_2 \cup F_3 \cup W \cup Z_i$. Thus, it is enough to show that Z_i can be ordered as $Z_i = \{z_1^i, \dots, z_{t_i}^i\}$ so that $N_G(z_{t_i}^i) \cap (F_2 \cup F_3 \cup W) \subseteq \dots \subseteq N_G(z_1^i) \cap (F_2 \cup F_3 \cup W)$. In view of Proposition 2.6, it suffices to show $F_2 \cup F_3 \cup W$ is a clique (since we already know that Z_i is a clique). But this follows immediately from Claim 4.5.2. ♦

Our proof is now complete: Claims 4.5.1-4.5.15 together guarantee that G is indeed a tent, and that $(A_0, A_1; B_0, B_1, B_2, B_3; C_1, C_2, C_3; F_2, F_3; W; Y; Z)$ is an associated tent partition. □

Theorem 4.6. *For any graph G , the following are equivalent:*

- G is $(2P_3, C_4, C_6, C_7)$ -free and contains an induced T_0 ;
- G is a tent.

Proof. This follows immediately from Proposition 4.4 and Lemma 4.5. $\square$

Corollary 4.7. *For any graph G , the following are equivalent:*

- (a) G is $(2P_3, C_4, C_6, C_7)$ -free, contains an induced T_0 , and contains no simplicial vertices;
- (b) G has exactly one nontrivial anticomponent, and this anticomponent is a thickening of T_0 or T_1 ;
- (c) G can be obtained from a thickening of T_0 or T_1 by possibly adding universal vertices to it.

Proof. Fix a graph G . By Propositions 2.1 and 4.1, (b) and (c) are equivalent.³¹ We will complete the proof by showing that (a) implies (c), and that (b) implies (a).

Suppose first that (a) holds, so that G is $(2P_3, C_4, C_6, C_7)$ -free, contains an induced T_0 , and contains no simplicial vertices. Since G is $(2P_3, C_4, C_6, C_7)$ -free, Theorem 4.6 guarantees that G is a tent. Let $(A_0, A_1; B_0, B_1, B_2, B_3; C_1, C_2, C_3; F_2, F_3; W; Y; Z)$ be a tent partition of the tent G . The result now readily follows from Proposition 4.3. Indeed, since G contains no simplicial vertices, Proposition 4.3(a-b) guarantees that $Y = Z = \emptyset$; but then Proposition 4.3(c) guarantees that all vertices in W are universal in G , and Proposition 4.3(d) guarantees that $G \setminus W$ is a thickening of T_0 or T_1 . Thus, (c) holds.

Suppose now that (b) holds. Then G has exactly one nontrivial anticomponent, and this anticomponent, call it Q , is a thickening of T_0 or T_1 . Since T_0 is an induced subgraph of T_1 , we see that Q (and therefore G as well) contains an induced T_0 . Further, by Proposition 4.3, Q is $(2P_3, C_4, C_6, C_7)$ -free and contains no simplicial vertices. But then Propositions 2.2 and 2.3 together guarantee that G is also $(2P_3, C_4, C_6, C_7)$ -free and contains no simplicial vertices. Thus, (a) holds. $\square$

5 The structure of $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or T_0

The following theorem gives a full structural description of $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or T_0 . Recall that the 7-saucer was defined in section 3.1, whereas the tent was defined in section 4.

Theorem 5.1. *For any graph G , the following are equivalent:*

- G is $(2P_3, C_4, C_6)$ -free and contains an induced C_7 or T_0 ;
- G is a 7-saucer or a tent.

Proof. This follows immediately from Theorems 3.9 and 4.6. $\square$

By combining Theorem 3.11 and Corollary 4.7, we obtain the following. (Recall that the family $\mathcal{M}$ was defined in subsection 3.2, and that graphs T_0 and T_1 are represented in Figure 1.2.)

Theorem 5.2. *For any graph G , the following are equivalent:*

³¹Indeed, by Proposition 4.1, any thickening of T_0 or T_1 is anticonnected (and it obviously contains at least two vertices), and so by Proposition 2.1, (b) and (c) are equivalent.

- (a) G is $(2P_3, C_4, C_6)$ -free, contains an induced C_7 or T_0 , and contains no simplicial vertices;
- (b) G has exactly one nontrivial anticomponent, and this anticomponent is a thickening of a graph in $\mathcal{M} \cup \{T_0, T_1\}$;
- (c) G can be obtained from a thickening of a graph in $\mathcal{M} \cup \{T_0, T_1\}$ by possibly adding universal vertices to it.

Proof. This follows immediately from Theorem 3.11 and Corollary 4.7. $\square$

6 Clique-width

The *clique-width* of a graph G , denoted by $\text{cwd}(G)$, is the minimum number of labels needed to construct G using the following four operations:

1. creation of a new vertex v with label i ;
2. disjoint union of two labeled graphs;
3. joining by an edge every vertex labeled i to every vertex labeled j (where $i \neq j$);
4. renaming label i to label j .

The main goal of this section is to prove Theorem 6.6, which states that $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 , and contain no simplicial vertices, have bounded clique-width. First, in subsection 6.1, we prove some technical results concerning the behavior of clique-width under “substitution.” Then, in subsection 6.2, we use Theorem 5.2 and the results of subsection 6.1 to prove Theorem 6.6.

6.1 Clique-width and substitution

Given graphs G and H on disjoint vertex sets, and given a vertex $v \in V(G)$, we say that a graph G^* is obtained by *substituting* H for v in G provided that the following hold:

- $V(G^*) = (V(G) \setminus \{v\}) \cup V(H)$;
- $G^*[V(H)] = H$;
- if $|V(G)| \geq 2$, then $G^* \setminus V(H) = G \setminus v$;³²
- for all $u \in V(G) \setminus \{v\}$, the following hold:
 - if $uv \in E(G)$, then u is complete to $V(H)$ in G^* ,
 - if $uv \notin E(G)$, then u is anticomplete to $V(H)$ in G^* .

Lemma 6.1. *Let G and H be graphs on disjoint vertex sets, let $v \in V(G)$, and let G^* be the graph obtained by substituting H for v in G . Then $\text{cwd}(G^*) = \max \{ \text{cwd}(G), \text{cwd}(H) \}$.*

Proof. Set $t := \max \{ \text{cwd}(G), \text{cwd}(H) \}$. Since H is an induced subgraph of G^* , and since G is isomorphic to an induced subgraph of G^* , it is clear that $\text{cwd}(G^*) \geq t$. It remains to show that $\text{cwd}(G^*) \leq t$. We proceed as follows. First, we independently construct the graph G using only labels $1, \dots, t$, as in the definition of clique-width; by symmetry, we may assume that the vertex v is assigned label 1 when it is first created during the construction of G . Next, we

³²If $|V(G)| = 1$, then $G \setminus v$ is not defined (since our graphs are assumed to be nonnull).

construct the graph H using only the labels $1, \dots, t$, as in the definition of clique-width, and then we rename all the labels used on H as 1. (So, at this point, all vertices of H have label 1.) We now follow the procedure that we originally followed in our construction of G , with one single modification: we do not create the vertex v , and instead, at the point when we would have created v , we take the disjoint union of the graph we have created and the graph H (with all vertices of H labeled 1).³³ This way, we construct the graph G^* using only labels $1, \dots, t$. Thus, $\text{cwd}(G^*) \leq t$. $\square$

Proposition 6.2. *For all positive integers k , we have that*

$$\text{cwd}(K_k) = \begin{cases} 1 & \text{if } k = 1, \\ 2 & \text{if } k \geq 2. \end{cases}$$

Proof. It is clear that $\text{cwd}(K_1) = 1$ and $\text{cwd}(K_2) = 2$. Moreover, it is clear that for all positive integers k , K_{k+1} can be obtained by substituting a copy of K_2 for a vertex of a copy of K_k . The result now follows from Lemma 6.1 via an easy induction on k . $\square$

Lemma 6.3. *Let G^* be a thickening of a graph G . Then $\text{cwd}(G) \leq \text{cwd}(G^*) \leq \max\{\text{cwd}(G), 2\}$. Consequently, if G has at least one edge, then $\text{cwd}(G^*) = \text{cwd}(G)$.*

Proof. Clearly, any graph that has at one edge has clique-width at least 2; therefore, the second statement follows from the first. Meanwhile, the first statement follows from Lemma 6.1 and Proposition 6.2, since G^* can be obtained from G by substituting complete graphs for the vertices of G . $\square$

Lemma 6.4. *If a graph G can be obtained from a graph Q by possibly adding universal vertices to Q , then $\text{cwd}(Q) \leq \text{cwd}(G) \leq \max\{\text{cwd}(Q), 2\}$.*

Proof. Fix graphs G and Q such that G can be obtained from Q by possibly adding universal vertices to Q . We must show that $\text{cwd}(Q) \leq \text{cwd}(G) \leq \max\{\text{cwd}(Q), 2\}$. We may assume that $G \neq Q$, for otherwise, the result is immediate. Since Q is an induced subgraph of G , it is clear that $\text{cwd}(Q) \leq \text{cwd}(G)$. It remains to show that $\text{cwd}(G) \leq \max\{\text{cwd}(Q), 2\}$.

Set $W := V(G) \setminus V(Q)$, and set $k := |W|$. Then G can be obtained from a copy of K_2 by first substituting Q for one vertex of the K_2 , and then substituting a copy of K_k for the other vertex of the K_2 . So, by Lemma 6.1, we have that $\text{cwd}(G) \leq \max\{\text{cwd}(K_2), \text{cwd}(Q), \text{cwd}(K_k)\}$. By Proposition 6.2, we have that $\text{cwd}(K_2) = 2$ and $\text{cwd}(K_k) \leq 2$, and we deduce that $\text{cwd}(G) \leq \max\{\text{cwd}(Q), 2\}$, which is what we needed to show. $\square$

Lemma 6.5. *Let G be a graph that can be obtained from a thickening of a graph H by possibly adding universal vertices to it. Then $\text{cwd}(G) \leq \max\{\text{cwd}(H), 2\} \leq \max\{|V(H)|, 2\}$.*

Proof. Let H^* be a thickening of H such that G can be obtained from H^* by possibly adding universal vertices to it. By Lemma 6.4, we know that $\text{cwd}(G) \leq \max\{\text{cwd}(H^*), 2\}$, and by Lemma 6.3, we know that $\text{cwd}(H^*) \leq \max\{\text{cwd}(H), 2\}$. So, $\text{cwd}(G) \leq \max\{\text{cwd}(H), 2\}$. Finally, it is clear that $\text{cwd}(H) \leq |V(H)|$ (indeed, we can create H by simply using a different label for each vertex of H), and the result follows. $\square$

³³After this point (i.e. the creation of H instead of v , with all vertices of H labeled 1), we follow exactly the same procedure as we did when we created G . (However, v never gets created, and instead, we have H .)

6.2 The clique-width of $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0

Theorem 6.6. *Let G be a $(2P_3, C_4, C_6)$ -free graph that contains an induced C_7 or an induced T_0 . Assume that G contains no simplicial vertices. Then $\text{cwd}(G) \leq 12$.*

Proof. By Theorem 5.2, we know that G can be obtained from a thickening of a graph in $\mathcal{M} \cup \{T_0, T_1\}$ by possibly adding universal vertices to it. Since every vertex in $\mathcal{M} \cup \{T_0, T_1\}$ has at most 12 vertices, Lemma 6.5 guarantees that $\text{cwd}(G) \leq 12$. $\square$

7 Some algorithmic considerations

$(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 are trivially recognizable in $O(n^7)$ time. Similarly, since T_0 has nine vertices, $(2P_3, C_4, C_6, C_7)$ -free graphs that contain an induced T_0 can be recognized in $O(n^9)$ time. Further, as stated in the introduction, GRAPH COLORING can be solved in polynomial time for graphs of bounded clique-width [25], and consequently, Theorem 6.6 guarantees that GRAPH COLORING can be solved in polynomial time for $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 . However, as pointed out in the introduction, coloring algorithms relying on clique-width are very slow (albeit polynomial).

Using our structural results, as well as various results from the literature, we can in fact recognize and color $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or an induced T_0 in only $O(n^3)$ time. We note, however, that the coloring algorithm relies on integer programming and is therefore not combinatorial.

7.1 Recognition

We sketch the algorithm for recognizing $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 ; the algorithm for recognizing $(2P_3, C_4, C_6, C_7)$ -free graphs that contain an induced T_0 is very similar, except that it relies on Theorem 4.6 and Corollary 4.7, rather than on Theorems 3.9 and 3.11.

So, let us suppose that we are given an input graph G for which we need to determine whether it is $(2P_3, C_4, C_6)$ -free with an induced C_7 . By Theorem 3.9, we just need to check if G is a 7-saucer. We start by computing a maximal sequence $v_1, \dots, v_t$ of vertices such that for all $i \in \{1, \dots, t\}$, the vertex v_i is simplicial in the graph $G \setminus \{v_1, \dots, v_{i-1}\}$. This can trivially be done in $O(n^4)$ time: we search for a simplicial vertex by examining the neighborhood of each vertex, and if we find a simplicial vertex, we delete it and repeat the process until no simplicial vertices remain. However, by using an algorithm described in the introduction of [17],³⁴ the sequence $v_1, \dots, v_t$ can in fact be found in only $O(n^3)$ time. Set $A := \{v_1, \dots, v_t\}$. If $V(G) = A$, then $v_1, \dots, v_t$ is a simplicial elimination ordering of G , and in this case, G is chordal (by [16]) and therefore does not belong to our class (because it does not contain an induced C_7). Let us now assume that $A \subsetneq V(G)$, and consider the graph $G' := G \setminus A$. We form the set W of all universal vertices of G' , and we form the graph $G'' := G' \setminus W$. Note that the twin relation in G'' is an equivalence relation,³⁵ and by [4, 26], the twin classes of G'' can be computed in $O(n^2)$ time.³⁶ Now, form the set X by taking exactly one vertex out of each twin class of G'' , and set $G_X := G[X]$. Clearly, G'' is a thickening of G_X . We then check if G_X is isomorphic to

³⁴Technically, we need a minor modification of this algorithm. See Lemma 2.5 of [20].

³⁵We technically mean the binary relation on $V(G'')$ such that vertices $x, y \in V(G'')$ are related if and only if $N_{G''}[x] = N_{G''}[y]$, that is, if and only if either $x = y$ or x, y are twins in G'' .

³⁶More precisely, an exercise from [26] asserts that such an algorithm exists, and a detailed proof can be found in [4].

any graph in $\mathcal{M}$; since all graphs in $\mathcal{M}$ have at most 12 vertices, this can be done in $O(1)$ time. If G_X is not isomorphic to any graph in $\mathcal{M}$, then Theorem 3.11 guarantees that G does not belong to our class, i.e. that G either is not $(2P_3, C_4, C_6)$ -free or does not contain an induced C_7 .³⁷ Let us now assume that G_X is isomorphic to a graph in $\mathcal{M}$. Then G' admits a special partition, and this special partition can easily be “reconstructed” using the twin classes of G_X and the set W . So, we now have a special partition $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W)$ of G' . It remains to examine A . If $A = \emptyset$, then G is a 7-saucer with an associated 7-saucer partition $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$, and we are done. So, assume that $A \neq \emptyset$. If A is not anticomplete to $X_0 \cup \dots \cup X_6$, or if there exists some $i \in \mathbb{Z}_7$ such that there is an edge between A and Y_i , and simultaneously, $Z_{i+2} \neq \emptyset$, then G is not a 7-saucer, and we are done.³⁸ Let us assume that A is anticomplete to $X_0 \cup \dots \cup X_6$, and that for all $i \in \mathbb{Z}_7$, either A is anticomplete to Y_i , or $Z_{i+2} = \emptyset$. We then compute the vertex sets of the components $A_1, \dots, A_\ell$ of A (using, for example, breadth-first-search), and we check if $A_1, \dots, A_\ell$ are all cliques. If not, then G is not a 7-saucer. So, let us assume that $A_1, \dots, A_\ell$ are indeed cliques. We then compute the degrees (in G) of all vertices in A , and for each $i \in \{1, \dots, \ell\}$, we order A_i as $A_i = \{a_1^i, \dots, a_{r_i}^i\}$ so that $N_G[a_{r_i}^i] \subseteq \dots \subseteq N_G[a_1^i]$, and we check if $N_G[a_{r_i}^i] \subseteq \dots \subseteq N_G[a_1^i]$. If we do indeed have that $N_G[a_{r_i}^i] \subseteq \dots \subseteq N_G[a_1^i]$ for all $i \in \mathbb{Z}_7$, then G is a 7-saucer with an associated 7-saucer partition $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$, and we are done. Otherwise, G is not a 7-saucer. Clearly, the running time of the algorithm is $O(n^3)$.

7.2 Coloring

Let us sketch an $O(n^3)$ coloring algorithm for $(2P_3, C_4, C_6)$ -free graphs that contain an induced C_7 or T_0 . We note that our algorithm is robust in the following sense: for any input graph G , we either produce an optimal coloring of G , or we correctly certify that the graph G is not in our class, i.e. that it is either not $(2P_3, C_4, C_6)$ -free, or it contains no induced C_7 and no induced T_0 .

Given an input graph G , we first compute a maximal sequence $v_1, \dots, v_t$ of vertices such that for all $i \in \{1, \dots, t\}$, the vertex v_i is simplicial in the graph $G \setminus \{v_1, \dots, v_{i-1}\}$; as explained in section 7.1, this can be done in $O(n^3)$ time. If $V(G) = \{v_1, \dots, v_t\}$, then $v_1, \dots, v_t$ is in fact a simplicial elimination ordering of G , and we can obtain an optimal coloring of G by simply coloring G greedily using the ordering $v_t, \dots, v_1$. So, assume that $\{v_1, \dots, v_t\} \subsetneq V(G)$, and set $G' := G \setminus \{v_1, \dots, v_t\}$. We then compute the set W of all universal vertices of G' , and we set $G'' := G \setminus W$. Once again as in subsection 7.1, we compute the twin classes of G'' in $O(n^2)$ time. Now, by Theorem 5.1, either G'' is a thickening of a graph in $\mathcal{M} \cup \{T_0, T_1\}$, or G does not belong to our class. But note that each graph in $\mathcal{M} \cup \{T_0, T_1\}$ has at most 12 vertices. We now simply check how many twin classes G'' has: if it has more than 12, then it is not a thickening of a graph in $\mathcal{M} \cup \{T_0, T_1\}$, and consequently, G does not belong to our class. So, let us assume that G'' has at most 12 twin classes. We can then use integer programming to compute an optimal coloring of G'' in $O(n + m)$ time, as explained in [18, 19].³⁹ We now finish coloring G in the obvious way: we assign all vertices of W distinct colors, not used on G'' , and then we complete our coloring by assigning colors to $v_t, \dots, v_1$ greedily. The total running time of our coloring algorithm is $O(n^3)$.

³⁷Technically, we are also using the fact that no graph in $\mathcal{M}$ contains a pair of twins.

³⁸The reader may possibly object that all that this proves is that $(X_0, \dots, X_6; Y_0, \dots, Y_6; Z_0, \dots, Z_6; W; A)$ is not a 7-saucer partition of G , but that it does not preclude the possibility that a different 7-saucer partition of G might exist. However, by examining the proof of Lemma 3.8, and keeping in mind that graphs that admit a special partition do not contain simplicial vertices (by Theorem 3.11), we see that this cannot possibly happen.

³⁹We do note that the $O(n + m)$ hides a large multiplicative constant.

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